

Convergence of Distributed Flooding and Its Application for Distributed Bayesian Filtering

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Abstract—Distributed flooding is a fundamental information sharing method to get network consensus via peer-to-peer communication. However, a unified consensus-oriented formulation of the algorithm and its convergence performance are not yet explicitly available in the literature. To fill this void in this paper, set-theoretic flooding rules are defined by encapsulating the information of interest in finite sets (one set per node), namely distributed set-theoretic information flooding (DSIF). This leads to a new type of consensus referred to as “collecting consensus,” which aims to ensure that all nodes get the same information. Convergence and optimality analyses are provided based on a consistent measure of the degree of consensus of the network. Compared with the prevailing averaging consensus, the proposed DSIF protocol benefits from avoiding repeated use of any information and offering the highest converging efficiency for network consensus while being exposed to increasing node-storage requirements against communication iterations and higher communication load. The protocol has been advocated for distributed nonlinear Bayesian filtering, where each node operates a separate particle filter, and the collecting consensus is pursued on the sensor data alone or jointly with intermediate local estimates. Simulations are provided in detail to demonstrate the theoretical findings.

Index Terms—Consensus, diffusion, distributed tracking, particle filter, sensor network.

I. INTRODUCTION

DISTRIBUTED computation has gained immense attention in the past decade, accompanying the rapid development and popularity of wireless sensor networks. In the successful networking operation, it is often of high interest that each node iteratively shares information with its intermediate neighbors (namely peer-to-peer communication) and consequently the entire network tends to reach a global alignment [1]/consensus [2]–[4] (to a certain degree). Compared to the centralized networking solutions based on a fusion center, distributed networking offers several advantages regarding scalability to adding or

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removing nodes, immunity to node failure, and dynamic adaptability to network topology changes.

However, there is a significant conflict between the degree of consensus (DoC) and communication requirement, as a higher DoC requires more communication, either more communicating iterations or higher communicating bandwidth which are limited by real time implementation and the communicating affordability of the nodes, respectively. Therefore, it is of paramount significance to seek a good balance for the trade-off so that the network achieves a satisfactory consensus in real-time and with affordable communication costs, which forms the majority of the research in the literature.

One of the most fundamental solutions for network information sharing is the flooding carried out in a distributed manner, by which all nodes synchronously broadcast their information to neighbors from the near to the distant. This protocol is well known in a few areas such as the communications [6]. Given that the network is strongly connected, it is able to achieve complete consensus (CC, i.e., all nodes have exactly the same set of information) after a certain number of iterations of peer-to-peer communication. This is quite appealing in theory but poses crucial challenges to the storage and communication requirement for large networks in practice. In fact, flooding has rarely been investigated in literature for distributed filtering (except for a few works, e.g., [7], [8]) even for small networks, regardless of its fast convergence (to be explicitly demonstrated in this paper) and ease of implementation. To note, there is one work that also refers to distributed flooding [9] which transmits the information of one single node over the network for routing. It is very different from the distributed protocol we consider here.

To date, a unified consensus-oriented formulation of the flooding algorithm and its convergence analysis are still missing in the literature. On one hand, there are cases in which DoC is required in the first priority while the sensor nodes have sufficient node-storage and communicating affordability to do so, for which flooding or even CC is simply preferable. On the other hand, instead of CC, it is more desired to perform flooding in a fewer, affordable, number of iterations for real time realization. What then can be expected and how can the number of iterations be properly determined?

To address the void and to answer the questions above, this paper aims to contribute from three aspects:

- 1) Formulate the flooding algorithm by encapsulating the information of interest in a finite set at each node and define set-theoretic flooding rules, namely *distributed set-theoretic information flooding* (DSIF), raising a new type of consensus termed *collecting consensus*.
- 2) Analysis the convergence and optimality of the DSIF protocol, based on a novel, consistent, metric of the DoC.

It is shown that DSIF enjoys the highest efficiency for network consensus among all distributed peer-to-peer communication schemes while suffering from heavier storage and communicational costs.

- 3) Show how the proposed DSIF scheme can be applied for and can benefit distributed Bayesian filtering, in which each node runs a separate particle filter (PF). Local PFs share sensor data alone or jointly with intermediate posteriors via DSIF. The latter usually needs to be parameterized in order to reduce the communication cost. The gain and loss to do so are analyzed and demonstrated in simulations.

The remainder of this paper is organized as follows. Notations and definitions regarding networking and three prevailing distributed information sharing protocols are given in Section II. As the main theoretical contribution, the collecting consensus-oriented DSIF is formulated in Section III with an analysis of its convergence and optimality. Section IV shows how to apply DSIF for distributed PF (DPF), with a brief literature review for DPF also given. Simulations are given in Section V and we conclude in Section VI.

II. NOTATION, DEFINITION AND BACKGROUND

A. Notation

The network topology is represented by a directed graph $G = (V, E)$ with the set of nodes $V = \{1, 2, \dots, N\}$ and the set of edges $E \subseteq V \times V$. In the directed graph, any edge is denoted by an ordered pair of nodes $(i, j) \in E$, which means node j is directly reachable from node i , where i is called the in-neighbor of j while j is the out-neighbor of i . For any $j \in V$, denote $N_j := \{i \in V | (i, j) \in E, i \neq j\}$, which is the set of all the in-neighbors of node j excluding node j itself. Undirected graph is a special type of directed graph where for any $(i, j) \in E$, we must have $(j, i) \in E$. If there exists a sequence of connected edges as follows

$$\{(i, \cdot), \dots, (\cdot, j)\} \subseteq E \quad (1)$$

This sequence of edges is called a path from node i to node j (denoted as $Path_{i, \dots, j}$) and node j is said to be “reachable” from node i . A digraph is said to be “strongly connected” (SC) if any node is reachable from all the other nodes, which is the digraph that can reach CC. For undirected graphs, SC is the same as connectivity [3].

The length of a path is given by the number of edges on that path. The length of the shortest path (perhaps, not unique) from node i to node j is called the distance from node i to node j , denoted as $D(i - j)$. Particularly, $D(i - i) = 0$ and if $(i, j) \in E$, $D(i - j) = 1$. We denote the set of all the nodes that are of distance $t \in \mathbb{N} = \{0, 1, 2, \dots\}$ to j as $N_j(t)$ and that of distance $t \in \mathbb{N}$ or smaller to j as $N_j(\leq t)$, namely $N_j(t) := \{n \in V | D(n - j) = t\}$, $N_j(\leq t) := \{n \in V | D(n - j) \leq t\}$. Obviously, we have $N_j(1) = N_j$, $N_j(0) = j$.

The largest distance between any two nodes, denoted as D_m , is called the diameter of the graph which is given by

$$D_m = \max_{i, j \in V} D(i - j) \quad (2)$$

Clearly, D_m only exist in SC networks. For any SC networks of at least two nodes, we have $D_m \in [1, N - 1]$, where the left bound corresponds to the fully connected network in which all nodes are in-neighbors of the others, while the right bound corresponds to the weakest connected network where all nodes are on a single chain in order, each having no more than two intermediate neighbors.

With particular regard to the distributed filtering problem, we assume that each node has independent abilities for: (i) filtering calculation, (ii) sensing to collect observations and (iii) communicating to neighbors. The communication is carried out in recursive iterations between neighboring nodes, each iteration consisting of sending no more than one data packet and receiving no more than one data packet. In addition, we need to clarify the following two definitions.

Definition 1 (real time communication): Communication is fully carried out between two successive observations and causes no sensor data missing or time-delay to the filter.

Definition 2 (communication bandwidth): The maximum size of the data packet that one node can send to or receive from its neighbor per communication.

B. Averaging/Maximum/Minimum Consensus

Here we assume that each node has a local scalar value, referred to as *state*, and it is of interest to compute the average of these values. An averaging consensus algorithm [2], [3] of zero communication time-delay is to reach an agreement regarding the state x_i each node has with local adapting dynamics $u_i(t)$, which can be written in discrete-time as

$$x_i(t + 1) = x_i(t) + u_i(t) \quad (3)$$

where $t \in \mathbb{N}$ denotes the communication iteration, $x_i(0)$ and $x_i(t)$ denote the initial and updated state of node i after iteration t , respectively.

In most of the time, the dynamics $u_i(t)$ is defined as

$$u_i(t) = \sum_{j \in N_i} \omega_{j \rightarrow i} (x_j(t) - x_i(t)) \quad (4)$$

where $\omega_{j \rightarrow i}$ is neighboring weight from node j to node i [2]–[4].

The complete convergence of *averaging consensus*: at iteration t states that, for any $i, j \in V$,

$$x_i(t) = x_j(t) \quad (5)$$

and asymptotically convergence (in the sense that $t \rightarrow \infty$)

$$\|x_i(t) - x_j(t)\| \leq \epsilon \quad (6)$$

where $\|x - y\|$ is a measure of the discrepancy between x and y , and ϵ is an error bound or margin [10].

In contrast to the average, one might be only interested in the maximum/minimum state, namely maximum/minimum consensus, which defines the iteration as

$$x_i(t + 1) = \max / \min \{x_i(t), \{x_j(t)\}_{j \in N_i}\} \quad (7)$$

The averaging consensus was well investigated in the community of control and systems [2]–[5]. One major concern is with the convergence, for which the spectral properties of the

graph Laplacian play a crucial role [2], [3], [10]-[12]. The second concern is raised by the information correlation/dependence among neighbors (especially when they own in part the same information). To account for this, fusion weights will be assigned to coordinate the nodes, e.g., covariance intersection [14],[15], which has inspired many strategies. In the absence of clear information about the correlation or relative quality of the information among nodes, two weighting methods have been proposed: one is to weight all nodes equally [16] and the other is to weight them according to the size of their neighborhood (named Metropolis weights [17]).

Particularly for the tracking problem, the sensors may update their observation frequently, preventing sufficient peer-to-peer communicating to get the network converge. This necessitates limiting the number of communicating iterations to gain a trade-off or compromise between a high DoC and little missing or time-delay of sensor data.

200 C. Gossip and Diffusion

To save communication, one alternative is to apply gossip to randomly choose fewer neighbors at each time (rather than to all neighbors) for averaging. It turns out that under mild conditions this process converges over time asymptotically [18]. Gossip based distributed filtering has been reported in, e.g., [19],[20]. However, gossip experiences the same problem as inefficient/repeated computations, for example, the same set (or largely similar set) of nodes repeatedly fusing their information at different points in time.

As another alternative, diffusion [22]–[24] performs only one iteration of peer-to-peer communication (i.e., the sensing and consensus time scales are the same), avoiding the problem of repeated use of any information. Based on it, the distributed Kalman filter (DKF) [23],[24] does not only share sensor data, but also local intermediate estimates through a diffusion update step. By this, the one-iteration-only communication is actually carried out on two types of data: the sensor data for the incremental update, and the estimates for the diffusion update. Hybrid fusion has been previously studied in the network using a fusion center, e.g., [21].

221 III. COLLECTING CONSENSUS AND DSIF

222 A. Collecting Consensus

In this paper, we are interested in an information sharing protocol that does not repeatedly use any information, and that will converge to CC in a definite number of iterations. By CC, we mean that all the nodes have exactly the same information of interest. Such a consensus model in which each node aims to collect information from all reachable nodes via the shortest paths is referred to as collecting consensus. Different to averaging consensus, the information from different nodes remains conditionally independent (or more precisely stated, unfused) until the end of the communication, which requires that nodes have sufficient storage allowance.

To perform collecting consensus, the information that needs to be communicated is encapsulated as a set, and the DSIF algorithm defines the information set dynamics (in contrast to

(3)) based on the union operation

$$I_i(t+1) = I_i(t) \cup u_i(t) \quad (8)$$

where $I_i(t)$ and $u_i(t)$ denote the existing and new incoming information set of node i at iteration $t \in \mathbb{N}$ respectively, $I_i(0)$ denotes the initial information set at node i with the size $|I_i(0)| = 1, \forall i \in V$, and $u_i(0)$ the initial dynamics.

As a result of CC, all nodes shall have exactly the same information set, i.e., $\forall i \in V, t \geq D_m$,

$$I_i(t) = \bigcup_{j \in V} I_j(0) \quad (9)$$

To this end, the DSIF algorithm consists of two stages:

- 1) In the starting iteration, each node collects information from all its in-neighbors

$$u_i(0) = \bigcup_{j \in N_i} I_j(0) \quad (10)$$

- 2) In the following iterations $t \in \mathbb{N}^+ = \{1, 2, \dots\}$, each node collects the new information that its in-neighbors have received at the preceding iteration

$$u_i(t) = \bigcup_{j \in N_i} \{I_j(t) \setminus I_j(t-1)\} \quad (11)$$

where $A \setminus B$ is the set difference of A and B , namely the set of all elements that are members of A but not of B and when $t \geq D_m$, we will actually have $u_i(t) = \emptyset$.

As shown in (11), the receiving neighbors will sort out the new received data, which they then transmit to their out-neighbors in the next iteration. In this process, the same information may be repeatedly received over edges, leading to information overuse and communication power waste, which is one defect of the naive flooding protocol, named implosion [6]. To avoid this, we define the set-theoretic information flooding rules in the following.

261 B. Set-Theoretic Flooding Rules

Rule 1 (data sending): Each node only sends to its out-neighbors the new information that has never been flooded before, and does so no more than once in each iteration.

Rule 2 (data accepting): Each node will not repeatedly take in the same information either from different in-neighbors or from the same node at different iterations, but only accept the information at its first arrival.

For both rules above, the data from each node shall be associated with a unique ID for distinguishing. Given these two rules respected, we will have $\forall i \in V, t \in \mathbb{N}$,

$$|I_i(t)| = |N_i(\leq t)| \quad (12)$$

To combat time-increasing storage requirement and communication load (when $t < D_m$), each element of data (often called a tuple) may be somehow compressed via e.g., dimension reduction [25] and polynomial encoding [13], under the premise that little or even no information would be lost and the data from different nodes remain conditionally independent.

It is worth noting that in the case of maximum or minimum consensus, there is neither a problem of information set-size

growing nor information overuse as the fusion result is always a single maximum or minimum value.

C. Convergence and Optimality of DSIF

To gain insights of the convergence of the proposed DSIF scheme, we need a metric to measure the DoC for collecting consensus, for which we propose a metric based on the size of the information set as follows.

Definition 3 (DoC): The DoC, denoted as C^o , of a network with N nodes, is defined as follows

$$C^o(t) = \frac{\sum_{i=1}^N |I_i(t)| - N}{N(N-1)} \quad (13)$$

where $t \in \mathbb{N}$ denotes the number of DSIF iterations that has been performed and in the following we limit it to $t \leq D_m$. On the DoC, we have the following theorem, which states the convergence property of the DSIF protocol.

Theorem 1: $0 \leq C^o(t_1) < C^o(t_2) \leq 1, \forall 0 \leq t_1 < t_2 \leq D_m$.

Proof: Before performing DSIF, each node has its original one unit of data, i.e., $|I_i(0)| = 1, \forall i \in V$. That gives $C^o(0) = 0$. After $t \geq D_m$ DSIF iterations, CC will be reached as all nodes will have the same information, i.e., $|I_i(t)| = N, \forall i \in V$. Furthermore, we have the following two straightforward Claims (for which we omit any proof):

Claim 1: As stated by (12), the size of the information set owned by sensor $i \in V$ will not be reduced during flooding except that data fusion or removal is taken, i.e., $|I_i(t_1)| \leq |I_i(t_2)|$ for any $0 \leq t_1 < t_2$.

Claim 2: Supposing two nodes g and q are of distance D_m (namely $D(g-q) = D_m$), for any $0 < t \leq D_m$, there must exist at least one node $j \in N_q(t)$ on $Path_{g \dots j \dots q}$ satisfying $D(j-q) = t$ whose information will arrive to node q exactly at iteration t and then, we have $|I_q(t-1)| + 1 \leq |I_q(t)|$.

From these two claims, we may conclude that $C^o(t_1) < C^o(t_2), \forall 0 \leq t_1 < t_2 \leq D_m$, to accomplish the proof. ■

Theorem 2: In the sense of DoC as given in (13), the proposed DSIF achieves the highest converging efficiency among all distributed peer-to-peer communication schemes.

Proof: From the definition of the distance between nodes and the DSIF peer-to-peer communication rules, we have two additional straightforward Claims:

Claim 3: All nodes whose information can reach node i in t iterations of peer-to-peer communication belong to $N_i(\leq t)$.

Claim 4: All nodes $q \in N_i(\leq t)$ will surely flood their information to node i in t DSIF iterations.

A combination of Claims 3 and 4 indicates that the DSIF will gain the largest possible $|I_i(t)|$ for any $i \in V$ and $t \in \mathbb{N}$ as claimed, which entails the converging optimality. ■

D. Trade-off between DoC and Number of Iterations

For a given network topology, the DoC is uniquely determined by the number of DSIF iterations. In turn, one can also determine the required number of iterations for a desired DoC, e.g., $T_c = 0.5$. That is, the DSIF stops at iteration t once

$$C^o(t) \geq T_c \quad (14)$$

Algorithm 1: DSIF operations at node i for DoC T_c .

INITIALIZATION:

1: $t \leftarrow 1; C^o(0) \leftarrow 0$

2: $I_i(t) = I_i(t-1) \cup \bigcup_{j \in N_i} I_j(t-1)$

RECURSIVE FLOODING ITERATION:

3: While $C^o(t) < T_c$

4: $t \leftarrow t + 1$

5: $I_i(t) \leftarrow I_i(t-1) \cup \bigcup_{j \in N_i} \{I_j(t-1) \setminus I_j(t-2)\}$

6: $C^o(t) \leftarrow \frac{\sum_{i=1}^N |I_i(t)| - N}{N(N-1)}$

7: End while

8: Return: $I_i(t)$

Algorithm 1 summarizes the communicating operations that need to be performed on node i for a given DoC T_c . 329 330

For a constant network of a known topology, the minimum number of iterations can be determined a priori by (14). However for time-varying dynamic networks, it needs to be calculated online via a consensus algorithm. To facilitate the use in time-varying networks without burdening any consensus procedures, we define the local DoC metric as follows: 331 332 333 334 335 336

Definition 4 (Local DoC): The DoC of node $i \in V$, denoted as C_i^o , after t DSIF iterations, is given as 337 338

$$C_i^o(t) = \frac{|I_i(t)| - 1}{N_k - 1} \quad (15)$$

where N_k is number of nodes in the network at time k , which needs to be estimated if unknown. From here we derive the following theorem as the local-node version of Theorem 1. 339 340 341

Theorem 3: $0 \leq C_i^o(t_1) \leq C_i^o(t_2) \leq 1, \forall 0 \leq t_1 < t_2 \leq D_m$ 342

Proof: This theorem states the same content as Claim 1. As a key difference to Theorem 1, the equality $C_i^o(t_1) = C_i^o(t_2)$ holds when and only when the number of nodes that are of distance t_1 to node $i \in V$ is the same as that of distance t_2 , i.e., $|N_i(t_1)| = |N_i(t_2)|$. ■ 343 344 345 346 347

Setting a threshold, e.g., $T_c = 0.5$ which can be the same or different for different nodes, on the desired local DoC, the consensus updating at node i may stop at iteration t once 348 349 350

$$C_i^o(t) \geq T_c \quad (16)$$

Furthermore, based on DoC we can define the convergence speed (CoS), either globally or locally, to measure the change of the size of the information set at each iteration, which also indicates the local real-time communication bandwidth. 351 352 353 354

Definition 5 (CoS): At iteration $t \in \mathbb{N}$, the global CoS, denoted as C^s , is defined as, 355 356

$$C^s(t) = C^o(t) - C^o(t-1) \quad (17)$$

and the local CoS of node $i \in V$ is defined as 357

$$C_i^s(t) = C_i^o(t) - C_i^o(t-1) \quad (18)$$

Theorem 4: $C^s(t) > 0, C_i^s(t) \geq 0, \forall 1 \leq t \leq D_m$ 358

Proof: The theorem is immediate from Theorems 1 and 3 as for any $1 \leq t \leq D_m, i \in V$, we have $C^s(t) = C^o(t) - C^o(t-1) > 0$ from Theorem 1, and $C_i^s(t) = C_i^o(t) - C_i^o(t-1) \geq 0$ from Theorem 3. ■ 359 360 361 362

Theorems 1, 3 and 4 entail an appealing property of the DSIF protocol which will not only converge definitively, but also 363 364

365 has a guaranteed converging speed that is globally positive and
 366 locally non-negative everywhere and at any iteration until CC
 367 is reached. We refer to this as *strong convergence*. It, however,
 368 also indicates a (non-negative) increasing storage requirement
 369 against communicating iterations. As an alternative to (16), we
 370 can build the predetermined threshold on the local CoS, e.g.,
 371 $T_s = 0.1$, then the minimum number of iterations t needs to
 372 satisfy

$$C_i^s(t) \leq T_s \quad (19)$$

373 But it is critical to note that we do not have any monotonicity
 374 on the CoS, e.g., $C^s(t_2) \leq C^s(t_1)$ or $C_i^s(t_2) \leq C_i^s(t_1)$ for $1 \leq$
 375 $t_1 < t_2 \leq D_m$. Therefore, the CoS at iteration t does not say
 376 anything of the CoS at iteration $t + 1$.

377 E. Comparison and Practical Consideration

378 Both metrics of DoC and CoS are clearly defined and easier
 379 to calculate than the one proposed for averaging consensus, e.g.,
 380 convergence rate [10]–[12], steady-state mean-square deviation
 381 [4] or disagreement vector [2], [3]. As indicated by Theorem 2,
 382 no peer-to-peer communication protocols converge faster than
 383 DSIF in terms of DoC. This superiority, however, is achieved
 384 at the expense of higher node storage requirements and heavier
 385 communication bandwidths. If the size of the data set at one
 386 node exceeds its communication bandwidth, multiple iterations
 387 will then be needed for that data set, otherwise data fusion is
 388 required to control the data size. In the former case, the required
 389 number of iterations will increase, while in the latter case the
 390 information completeness or independence may not be kept.
 391 However, we will not address this issue further here, which is
 392 quite problem dependent. In brief, we have the following remark
 393 on the respective advantages of averaging consensus, diffusion
 394 and collecting consensus.

395 *Remark 1:* The averaging consensus takes the lowest com-
 396 municating bandwidth (always one unit of data) but more iter-
 397 ations to reach any DoCs while the diffusion severely limits
 398 the number of iterations (to one only) which may insufficiently
 399 use the communication affordability (i.e., more iterations are
 400 actually allowed in real time communication). In contrast, the
 401 proposed DSIF protocol aims to get the best possible consensus
 402 in an real-time-allowed number of iterations, which is therefore
 403 particularly suited to small and moderate networks for which
 404 the nodes have sufficient storage and communicating power. A
 405 means to facilitate its use in large networks is to selectively
 406 apply data fusion such as averaging in every several flooding
 407 iterations in order to control the data-set size. This will lead to
 408 a hybrid protocol that iterates between flooding and averaging
 409 consensus, to gain a balance between benefiting from high com-
 410 munication efficiency and suffering from information overuse
 411 and slower convergence.

412 IV. DISTRIBUTED BAYESIAN FILTERING USING DSIF

413 A. State-of-the-art DPF Protocols

414 Before presenting our DPF framework based on DSIF, a brief
 415 revisit of the PF algorithm and existing DPF protocols is given

below. Suppose that at time k , the local (marginal) posterior at
 sensor i is represented by a local PF

$$p(\mathbf{x}_k | \mathbf{z}_{i,1:k}) \approx \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} \delta(\mathbf{x}_k - \mathbf{x}_{i,k}^{(m)}) \quad (20)$$

where $\delta(\mathbf{x} - \mathbf{y})$ is the Dirac delta impulse, which equals to
 one if $\mathbf{x} = \mathbf{y}$ and to zero otherwise, \mathbf{x}_k is the true state vector,
 $\mathbf{z}_{i,1:k}$ is the observation serial, $\mathbf{x}_{i,k}^{(m)}$ and $w_{i,k}^{(m)}$ are the state and
 normalized weight of the m th particle respectively, $M_{i,k}$ is the
 total number of particles at filtering time k .

The essence of the PF is to assess how well each particle
 conforms to the state model and explains the observations, using
 this assessment to generate a weighted sample approximation
 to the Bayesian posterior, and thereby form sub-optimal state
 estimates. Given local measurement $\mathbf{z}_{i,k}$, $i \in V$, the weights of
 the particles are evaluated over time based on the sequential
 importance sampling (SIS) principle as

$$w_{i,k}^{(m)} \propto w_{i,k-1}^{(m)} \frac{p(\mathbf{z}_{i,k} | \mathbf{x}_{i,k}^{(m)}) p(\mathbf{x}_{i,k}^{(m)} | \mathbf{x}_{i,k-1}^{(m)})}{\pi(\mathbf{x}_{i,k}^{(m)} | \mathbf{x}_{i,k-1}^{(m)}, \mathbf{z}_{i,1:k})} \quad (21)$$

where $\pi(\cdot)$ is a proposal to generate particles, and in general
 its design shall take into account both the newest measure-
 ment $z_{i,k}$ and the prior in order to best match the posterior; see
 e.g. [19], [28], [31]. The use of the observation in the sampling
 proposal design is particularly helpful (and even necessary for
 avoiding sample degeneracy) when the observation is very ac-
 curate. However, caution should be exercised here since the
 repeated use of the observation (both for proposal design and
 in likelihood calculation) may not benefit the filter when the
 observation suffers from significant noise [42].

In addition to SIS, resampling is usually required to reduce
 the weight variance when it exceeds a certain threshold, so that
 all particles will have equal or approximate weights while the
 posterior distribution can be the best maintained [31], [32]. This
 has often been referred to as sampling importance resampling
 (SIR), which is the core of the majority of existing PFs. We
 assume the reader is familiar with the centralized PF and so
 limit ourselves hereafter to the distributed implementation, in
 which local nodes carry out PF calculations in parallel and
 meanwhile share information with their neighbors to assist their
 filters. For this, a variety of information sharing protocols have
 been proposed, which can be classified as follows:

- 1) *Sequential information passing:* Information transmits in
 a sequential, predefined manner from a node to one of
 its neighboring nodes via a cyclic path until the entire
 network is traversed [43]. The sequential realm is sensitive
 to the mobility and failure of nodes/edges and is time-
 consuming.
- 2) *Flooding:* As addressed, the flooding protocol provides
 the fastest albeit communication-intensive way to spread
 information over the network [7], [8], but, neither any clue
 to determine the number of communication iterations in
 order to compromise real time realization and DoC nor
 any convergence results has been shown.

3) *Averaging consensus*. There is a large body of work concerning averaging consensus-based distributed filtering. The data transmitted between neighboring nodes can be posterior statistics in the form of Gaussian component [33] /GM [29]–[30] or generalized probability densities [36]–[37], likelihood [26]–[28], particle set [34]–[35] or raw observations [38]. Excellent surveys are also available such as a taxonomy of DPFs [39], a comparison of several belief consensus algorithms [40] and a recent survey of convergence and error propagation of DPFs [28]. In summary, complete information sharing affords better accuracy but has higher communication requirements, such as [34]–[35] that exchange all particles. Parameter approximation [26]–[33] or random gossip [19]–[20] can significantly reduce the communication cost, but may lead to a deterioration in the filter performance.

4) *Diffusion*: The diffusion scheme addressed in Section II.C also provides a competitive alternative to the averaging consensus for DPF [7]–[8], [41].

We note that the sensor data can be either simple (e.g., range, bearing) or complex (e.g., image data). To avoid distracting from the key contribution of this paper on collecting consensus and the DSIF protocol, we only consider the former case for simplicity. For the latter case, one may consider compressing the sensor data, e.g., [25]–[26], [12]–[13] or transmitting the low-dimensional likelihood for replacement [26]–[27]. At the current stage, we have not considered complicated network issues such as communication constraints, e.g., [45]–[46], and asynchronous sensing, e.g., [47]–[48]. However, we note all of these issues are valuable to be investigated on the base of the proposed DSIF protocol.

B. DSIF on Sensor Data and on Local Posterior

In the proposed DPF framework, the DSIF scheme will be applied on the sensor data alone or jointly on local posteriors. In the latter, we propose parameterizing the posterior to save communication. Since a vast number of random numbers are required by the PF, it is communication intensive to run consensus on them, and it is not our intention to do so.

First, DSIF is implemented on the sensor data including the target-observations (and uncertainties) associated with the sensor ID, all as one unit. To note, the sensor position is often required for likelihood calculation and therefore can serve as the unique sensor ID for distinguishing. Then, the resultant consensus on sensor data with sensor profiles given a priori, is equivalent to collecting consensus on the likelihood which is required for PF updating. A likelihood function contains the information of both the sensor data and the sensor profile in a more compact manner. But for simplicity of understanding, we keep addressing consensus on sensor data.

The filtering posteriors obtained at different nodes, referred to as local posteriors, will be different, even if CC is reached on sensor data over the network where the difference attributes to the different random numbers. If DoC is low on sensor data, the difference between local posteriors will be relatively significant. As such, we may apply the second DSIF scheme to fuse local posteriors among neighbors as well as to get the local LMS

(least mean squares) estimate; we refer to this step as diffusion, in parallel to [24]. By this, each node aims to improve their local estimate with regard to their neighbors' posterior. However, parameter approximation of local posteriors, typically via Gaussian or GM approximation, is needed (otherwise massive communication will be triggered if the complete posterior is communicated by transmitting the entire particle set), which will in turn introduce approximation errors to the posterior. This trade-off is much problem-dependent and will determine whether the second DSIF is worthwhile.

The operations that need to be conducted on each sensor in the proposed distributed PF is summarized in Algorithm 2. In it, steps 1-a and 1-b are independent of each other and therefore can be carried out in either order or in parallel. Sensor data DSIF and posterior DSIF have been implemented t_1 and t_2 iterations respectively, where t_1 and t_2 are not necessarily equal but are determined for respective desired or the largest affordable DoCs as addressed in Section III. They show complementary features and resemble the Incremental and Diffusion updates of the diffusion-based DKF [24]. But, there are obvious differences:

- 1) Our framework is developed for nonlinear models which releases the requirement of linear system functions and even Gaussian assumption of the posterior;
- 2) Our consensus protocol does not limit information sharing between neighbors to one iteration only but instead, the DoC will be pursued as much as the real time communication allows;
- 3) Our diffusion update (Step 5 in Algorithm 2) is an optional step, which is advocated for re-setting local posteriors only when local posteriors are significantly different (as a consequence of a low DoC on the sensor data achieved in the first DSIF implementation). When the difference between local posteriors is insignificant (because of a high DoC achieved on the sensor data), there will be less need to further fuse them and so it may be better not to diffuse local posteriors since the errors introduced due to parameterization can be more significant than the benefit. This is a critical point. We will demonstrate this in detail through simulations in Section V. In addition, we provide two easy-to-implement diffusion choices.
- 4) We point out that the proposed two DSIF procedures can be performed jointly, although this may not reduce the communication load and the storage requirement in total; see the following Remark 2.

Remark 2: Two DSIF implementations regarding the sensor data and the local filter estimates form the starting step and the end step of each filtering iteration, respectively. In the time series, they are adjacent. Therefore, they may be combined in one joint consensus scheme at some stages (which however does not necessarily indicate that $t_1 = t_2$), i.e., the local estimates obtained at filtering time k can be combined with sensor data received at time $k + 1$ as one unit of data, both sharing the same node ID for DSIF. Then, the initial information set at node $i \in V$ can be defined as

$$I_{i,k}(0) := \{\hat{\mathbf{x}}_{i,k-1}, P_{i,k-1}\} \cup \mathbf{z}_{i,k} \quad (22)$$

Algorithm 2: Distributed PF calculation executed on node i .

Step 1-a Filter prediction: Propagate the particles $\mathbf{x}_{i,k-1}^{(m)}$ to $\mathbf{x}_{i,k}^{(m)}$ for all $m = 1, \dots, M_{i,k}$, according to the proposal function $\pi(\mathbf{x}_{i,k}^{(m)} | \mathbf{x}_{i,k-1}^{(m)}, \mathbf{z}_{i,1:k})$. At $k = 0$, particles are sampled from an initial proposal π_0 instead for filter initialization.

Step 1-b 1st DSIF: Perform t_1 DSIF iterations on sensor data as given in Algorithm 1, resulting in a combined measurement set $\mathbf{Z}_{i,k} = \{\mathbf{z}_{j,k}\}_{j \in N_i(\leq t_1)}$. This step is carried out whenever new measurements become available.

Step 2 Filter updating: Re-weight all particles via (21) ($\mathbf{z}_{i,t}$ therein shall be replaced by $\mathbf{Z}_{i,k}$ obtained in the 1st DSIF) and then normalize them as follows

$$w_{i,k}^{(m)} \leftarrow w_{i,k}^{(m)} / \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} \quad (23)$$

Step 3 Estimate extraction: Extract local estimate $\hat{\mathbf{x}}_{i,k}$ and calculate their covariance $P_{i,k}$ from the local random measure $\chi_{i,k} = \{\mathbf{x}_{i,k}^{(m)}, w_{i,k}^{(m)}\}_{m=1,2,\dots,M_{i,k}}$ as follows

$$\hat{\mathbf{x}}_{i,k} = \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} \mathbf{x}_{i,k}^{(m)} \quad (24)$$

$$P_{i,k} = \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} \left(\mathbf{x}_{i,k}^{(m)} - \hat{\mathbf{x}}_{i,k} \right) \left(\mathbf{x}_{i,k}^{(m)} - \hat{\mathbf{x}}_{i,k} \right)^T \quad (25)$$

Step 4 2nd DSIF: Perform t_2 DSIF on local estimates obtained in Step 3, resulting in a set of intermediate estimates $\{\hat{\mathbf{x}}_{j,k}, P_{j,k}\}_{j \in N_i(\leq t_2)}$, which will be fused in the LMS sense as follows

$$\hat{\mathbf{x}}_{i,k}^{\text{LMS}} = \sum_{j \in N_i(\leq t_2)} \hat{\mathbf{x}}_{j,k} P_{j,k}^{-1} \quad (26)$$

$$P_{i,k}^{\text{LMS}} = \left(\sum_{j \in N_i(\leq t_2)} P_{j,k}^{-1} \right)^{-1} \quad (27)$$

This also offers the local filter output at sensor i .

Step 5 Diffusion: As an option, the shared filter estimates given in Step 4 can be used to re-set the local PF posterior $\chi_{i,k}$. If so, there are two choices (the second is expected to have a higher approximation accuracy than the first).

1) Re-set $\chi_{i,k}$ as the LMS fused Gaussian distribution, as is done in the Gaussian PF [49].

$$\chi_{i,k} \leftarrow \mathcal{N}(\hat{\mathbf{x}}_{i,k}^{\text{LMS}}, P_{i,k}^{\text{LMS}}) \quad (28)$$

2) Re-set $\chi_{i,k}$ as the shared GM before performing LMS fusion, as is done in the Gaussian sum PF [50]

$$\chi_{i,k} \leftarrow \sum_{j \in N_i(\leq t_2)} \mathcal{N}(\hat{\mathbf{x}}_{j,k}, P_{j,k}) \quad (29)$$

Step 6 Resampling: Sample from the updated particle set [32] if the variance of weights exceeds a specified threshold and if Step 5 is not applied. If Step 5 is applied, sample from the diffused Gaussian or GM distribution $\chi_{i,k}$ given by Step 5 to generate a new particle set. Update $k \leftarrow k - 1$ and go to the next filtering iteration.

V. SIMULATIONS

575

In this section, we consider tracking a target that moves in the $x - y$ plane by using the proposed DPF based on a constant sensor network earlier appeared in [39] as given in Fig. 1. The network has totally 10 sensors and a diameter $D_m = 4$. The simulation models and parameters are the same to [27]. In specific, we have the initial state as $\mathbf{x}_0 = [4, 0.5, 4, 0.5]^T$. The Markov transition model that governs the target movement of nearly constant velocity is given by

$$\mathbf{x}_k = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \mathbf{x}_{k-1} + \begin{bmatrix} 0.5 & 0 \\ 1 & 0 \\ 0 & 0.5 \\ 0 & 1 \end{bmatrix} \mathbf{u}_k \quad (30)$$

where $\mathbf{x}_k = [p_{x,k}, \dot{p}_{x,k}, p_{y,k}, \dot{p}_{y,k}]^T$, $[p_{x,k}, p_{y,k}]^T$ gives the position and $[\dot{p}_{x,k}, \dot{p}_{y,k}]^T$ the velocity, $\mathbf{u}_k \sim \mathcal{N}(\mathbf{0}_2, 0.00035\mathbf{I}_2)$.

The target emits an acoustic or radio signal with a known constant transmit power P_t that can be received by all sensors independently, i.e., the scalar measurement function of sensor i located at $[s_{i,x}, s_{i,y}]^T$ about target \mathbf{x}_k is

$$z_{i,k} = \frac{\alpha P_t}{\| [p_{x,k}, p_{y,k}]^T - [s_{i,x}, s_{i,y}]^T \|^{\gamma}} + v_k \quad (31)$$

where α is a constant that depends on several factors such as fast and slow fading, and gains in the transmitter and receiver antennas, γ is the path loss exponent [44], and $v_k \sim \mathcal{N}(0, \sigma_v^2)$ is the measurement noise. In parallel to [27], we set simply $\alpha P_t = 10$, $\gamma = 1$, $\sigma_v^2 = 0.001$.

When multiple synchronous observations are available, the weight of particles is updated by multiplying the likelihoods given by each available measurement. That is,

$$p(Z_{i,k} | \mathbf{x}_{i,k}^{(m)}) = \prod_{j \in N_i(\leq t)} p(z_{j,k} | \mathbf{x}_{i,k}^{(m)}) \quad (32)$$

where $Z_{i,k} = \{z_{i,k}\}_{j \in N_i(\leq t)}$ is the measurement set at sensor i gained in the first DSIF procedure of total t iterations.

For any sensor $i \in V$, the necessary and sufficient number of iterations, denoted as $D_{m,i}$, to receive the information from all the other sensors can be given by

$$D_{m,i} := \max_{j \in V} D(j - i) \quad (33)$$

We design three groups of simulations in the following three subsections that use the same ground truths to evaluate or compare the following five PF protocols, where the first three are distributed while the last two are centralized. All PFs use the same number of particles ($M = 1000$).

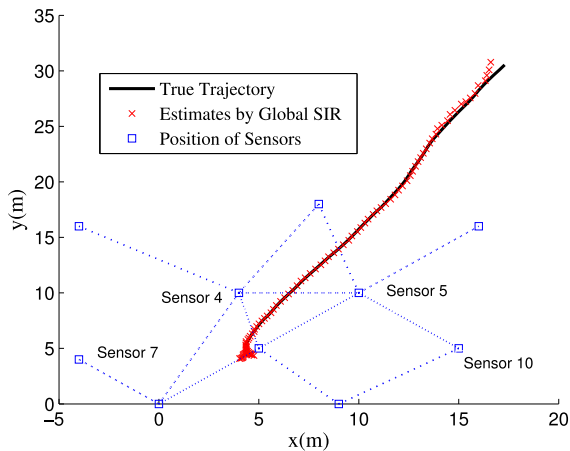


Fig. 1. The topology of the sensor network, the target trajectory and its estimate given a by a global SIR filter in one trial.

- 1) *C-SIR*: we apply DSIF only on the sensor data, named Consensus without Diffusion (i.e., Steps 4 and 5 are not applied in Algorithm 2). In this case, each local PF is a SIR filter that is free of any Gaussian assumption;
- 2) *CD-GMPF*: we apply DSIF on both sensor data and local estimates named Consensus with Diffusion (i.e. Steps 4 and 5 are applied in Algorithm 2). In this case, each local PF is a Gaussian sum PF that applies (29) for posterior approximation and fusion;
- 3) *L-C-SIR*: the Likelihood Consensus-based SIR filter [27] can be viewed as a special case of our C-SIR filter that applies sensor data averaging consensus (for likelihood multiplying) at each iteration. For fast converging, the Metropolis weights strategy [52] is employed for averaging in the L-C-SIR filter;
- 4) *Local-SIR/GMPF*: local SIR filter or GMPF that does not communicate with each other at all;
- 5) *Global-SIR*: a centralized SIR filter that is able to access all sensor observations at all times.

To mitigate the problem of sample impoverishment that is often caused by resampling in the SIR filters, the minimum-sampling-variance resampling [32] is applied when and only when the effective sample size is smaller than $M/2$ and if applied, a roughening noise that is equivalent to half of \mathbf{u}_k will be used [31].

To measure the filtering accuracy, we calculate the root mean square error (RMSE) on both the position estimate and the velocity estimate, respectively, as follows

$$\text{RMSE}_{\text{pos}_k} = \sqrt{\frac{1}{C} \sum_{c=1}^C (x_{k,c} - \hat{x}_{k,c})^2 + (y_{k,c} - \hat{y}_{k,c})^2} \quad (34)$$

$$\text{RMSE}_{\text{vel}_k} = \sqrt{\frac{1}{C} \sum_{c=1}^C (\dot{x}_{k,c} - \hat{\dot{x}}_{k,c})^2 + (\dot{y}_{k,c} - \hat{\dot{y}}_{k,c})^2} \quad (35)$$

where $[\hat{x}_{k,c}, \hat{y}_{k,c}]^T$ and $[\hat{\dot{x}}_{k,c}, \hat{\dot{y}}_{k,c}]^T$ are the position-estimate and velocity-estimate given at filtering time k in trial c , respectively, and $C = 20$ is the total number of MC trials. Further, the

TABLE I
DoC ACHIEVED AT EACH DSIF ITERATION (LOCAL AND GLOBAL)

	Sensor 4	Sensor 5	Sensor 7	Sensor 10	Global
$t = 0$	0	0	0	0	0
$t = 1$	5/9	5/9	1/9	2/9	26/90
$t = 2$	1	8/9	3/9	6/9	61/90
$t = 3$	1	1	7/9	8/9	86/90
$t = 4$	1	1	1	1	1

average position RMSE is defined as the mean of $\text{RMSE}_{\text{pos}_k}$ over the entire simulation period of 100 filtering iterations. In each trial, the ground truth is independently generated (for generality). In all trials, the prior distribution of the particle set is initialized around the true state as $\mathcal{N}(x_0, P_0)$, with $P_0 = \text{diag}[2, 0.001, 2, 0.001]^T$.

In particular, we will assess the filter performance at four representative sensors, marked in Fig. 1 as sensors 4, 5, 7 and 10. For them, we have $D_{m,4} = 2, D_{m,5} = 3, D_{m,7} = 4, D_{m,10} = 4$. This means that sensor 4 will achieve CC first (after 2 iterations) while sensors 7 and 10 will be the last (after 4 iterations). For different numbers of DSIF iterations, the global and local DoCs are given in Table I. Particularly, for $t = 1$, we have the global DoC determined as

$$C^o(1) = \frac{|E|}{N(N-1)} \quad (36)$$

where $|E|$ is the number of edges; (a, b) and (b, a) are counted as two different edges.

A. Consensus without Diffusion

In this case, each sensor operates a separate SIR filter. Sensors are assumed conditionally independent and use different random numbers. The posteriors obtained by sensors will be different from both each other and the global/local PF, even given that they all reach CC on sensor data.

For different numbers of DSIF iterations from 0 (no consensus at all) to 4 (D_m), the RMSEs of the position and velocity estimation of local C-SIR filters and the global SIR filter are given in Fig. 2.(a)-(e) respectively, corresponding to different DoCs. The average RMSEs over 100 filtering steps against the number of DSIF iterations are given in Fig. 2.(f). The results clearly demonstrate that:

- 1) A single passive sensor is not capable of delivering good tracking in this problem as the RMSEs given by local PFs are much higher than that provided by DPFs; this necessitates the collaboration of multiple geographically dispersed sensors;
- 2) The more informative sensor data used, the better the filter performance;
- 3) The larger DoC, the closer the local PF performance to the centralized PF, i.e., local filters converge to the global filter against iterations as the DoC increases;
- 4) Once CC is reached, the performance of the local PF is very close to that of the centralized PF (with regard to

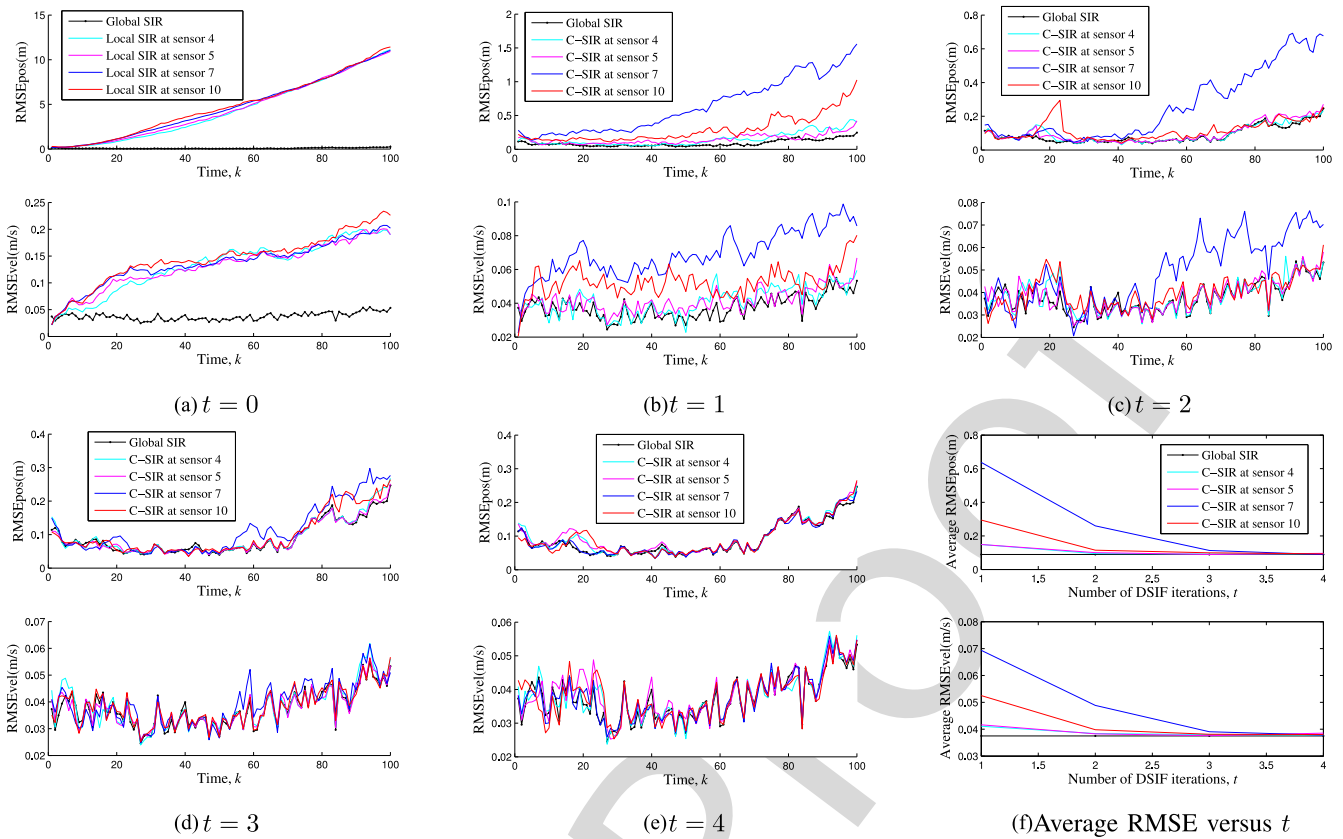


Fig. 2. Position and velocity RMSE of C-SIR filters with different numbers of DSIF iterations, comparing with the global SIR filter.

680 both position and velocity) but still not the same, since
 681 different random numbers are used.

682 Based on the measure of DoC, we are able to approximately
 683 determine how much information divergence different nodes
 684 will have and what payoff can be expected if one more or one
 685 less iteration of peer-to-peer communication is employed. For
 686 example, when the number of iterations is $t = 3$, the global
 687 DoC is as high as 86/90, close to 1, which agrees with the slight
 688 difference between Fig. 3(a) and (b). This is a valuable part of
 689 the metric of DoC.

690 B. Consensus with Diffusion

691 In this case, each sensor runs a separate GMPF. Collecting
 692 consensus are applied on both the sensor data and intermediate
 693 estimates jointly in a single DSIF procedure (and set $t_1 = t_2$).
 694 Because of the GM diffusion of intermediate estimates, the
 695 local fused estimates are expected to be closer to each other.
 696 If CC is reached, they shall be exactly the same. In parallel to
 697 the last simulation, different numbers of DSIF iterations from
 698 0 to 4 are employed to the CD-GMPFs, which are compared with
 699 the (centralized) global SIR PF in Fig. 3(a)–(e) respectively.
 700 The average RMSEs of these filters against the number of DSIF
 701 iterations are given in Fig. 3(f).

702 We use the same ground truth (20 MC trials) regarding the
 703 trajectories and sensor observations as the last simulation. Com-
 704 pared to the last simulation, we can find that

705 1) A single passive sensor can still hardly work well when
 706 the local SIR filters are replaced by local GMPFs;

- 707 2) Given the same number of DSIF iterations $t = 1, 2$, CD-
 708 GMPFs perform much better than C-SIR and are much
 709 closer to each other; this is because of the second DSIF
 710 scheme on the posteriors over the network which enhances
 711 the consensus to improve local estimates;
- 712 3) Given $t = 3, 4$ iterations, the local CD-GMPFs perform
 713 almost the same but different to the global SIR filter;
- 714 4) Given CC achieved, the RMSEs of all local GMPFs are
 715 exactly the same but are inferior to the global SIR fil-
 716 ter, especially at the later stage in this tracking example.
 717 Analysis and discussion will be given next.

718 C. Comparison and Discussion

719 Finally, we compare both types of DSIF-based DPFs with the
 720 L-C-SIR filter [27], [39]. The key difference of the likelihood
 721 consensus to DSIF is that each node fuses information interme-
 722 diately after receiving them and therefore the communication
 723 cost is lower, but it is exposed to repeated use of information
 724 and slower convergence.

725 First, for $t = 4$, the average (over all nodes) position RMSEs
 726 of the C-SIR, CD-GMPF and the L-C-SIR filters are given
 727 in Fig. 4. It shows that the C-SIR filter achieves the closest
 728 performance to that of the centralized filter. We further calculate
 729 the mean of these average RMSEs for $t = 0$ to 8 and the results
 730 are given in Fig. 5. It shows that these consensus protocols can
 731 all significantly improve the filter performance as compared to
 732 the local filter that applies no consensus and converges against
 733 communication iterations.

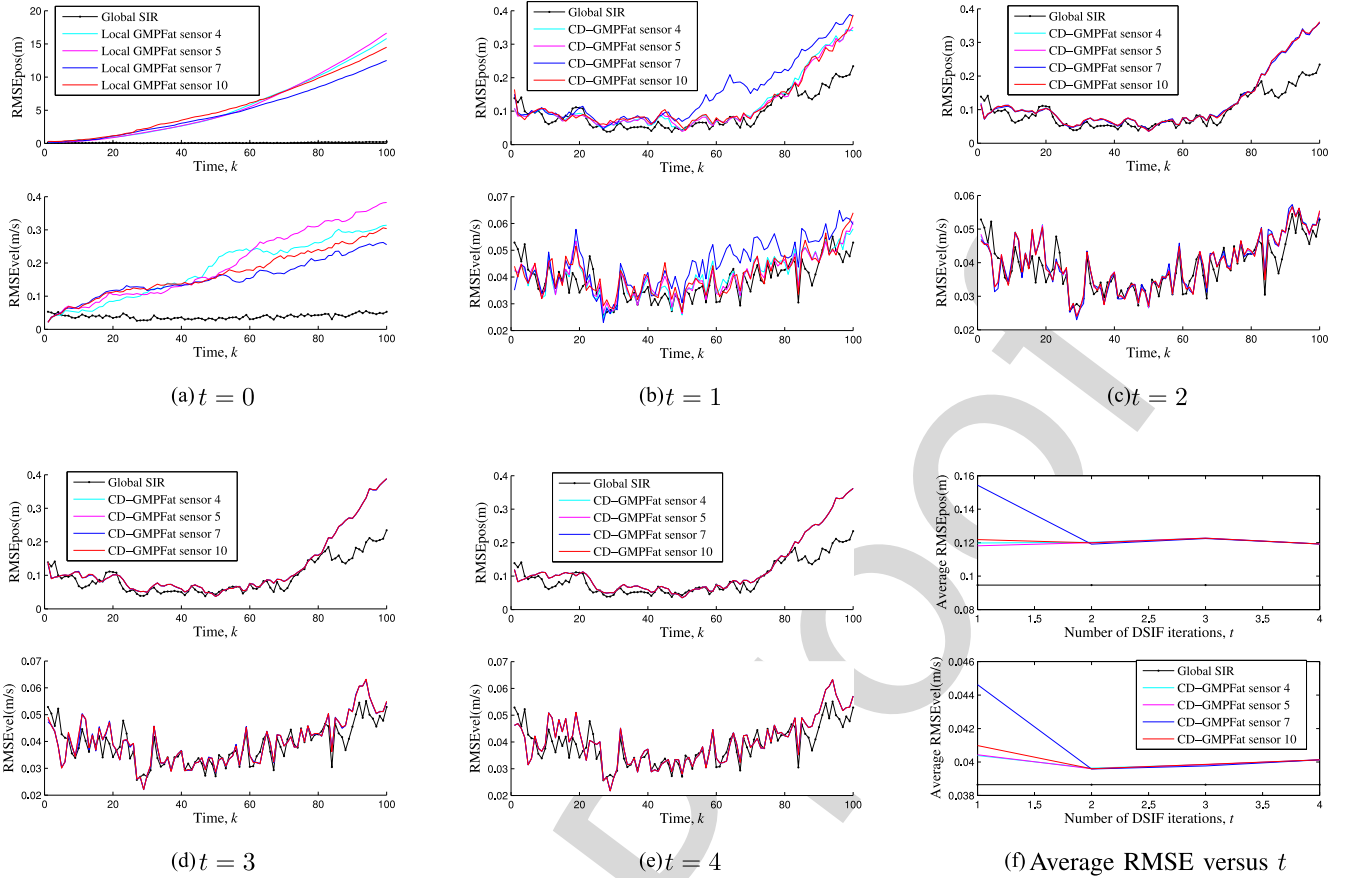


Fig. 3. Position and velocity RMSE of CD-GMPFs with different numbers of DSIF iterations, comparing to the global SIR filter.

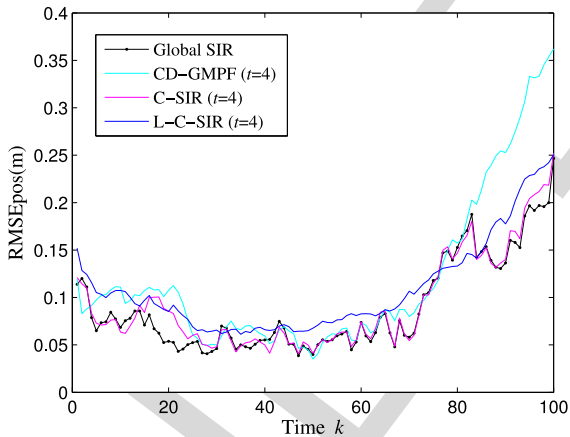


Fig. 4. Position RMSE of different DPFs applying 4 iterations of peer-to-peer communication.

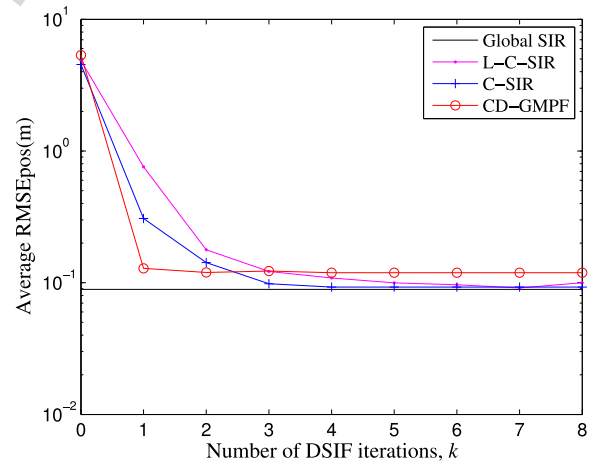


Fig. 5. Average position RMSE of different DPFs over 100 filtering steps against the number of peer-to-peer communication iterations.

734 Furthermore, we have the following observations, which
735 show more insights of these three types of DPFs:

- 736 1) DSIF based C-SIR and CD-GMPF converge faster than
737 the averaging consensus-based L-C-SIR filter at the ex-
738 pense of higher communication cost. CD-GMPF con-
739 verges the fastest but it suffers from a larger RMSE at
740 the end, all due to its diffusion step that shares infor-
741 mation among nodes more thoroughly than without diffu-
742 sion but also introduce errors;

- 2) For a relatively small number of iterations that correspond
743 to a low DoC on observation (which may lead to a large
744 discrepancy between local nodes' posteriors), the C-SIR
745 filter is inferior to the CD-GMPF, as shown in Fig. 5 (also
746 told by comparing between Figs. 2(f) and 3(f)). In this
747 case, the diffusion update leads to earlier convergence
748 and better performance for the filter. This is in line with
749 the findings reported in [23];
750

3) For a large number of iterations that correspond to a high DoC on observation and consequently on posterior (leaving little space to benefit from posterior fusion), the diffusion update of the CD-GMPF is not so preferable; instead, the GM approximation error caused in the diffusion might be more significant than the benefit it can offer, resulting in an overall filter degradation. We must note that if the whole particle sets are transmitted for diffusion without any approximations, and also the dependence between the posteriors are accounted for properly in the diffusion update, it shall always be beneficial in theory regardless of the much greater cost in communication and local fusion calculation.

These results confirm our theoretical prediction and demonstrate further that, both approximation and data fusion during communication can be either beneficial or counterproductive. Generally speaking, parametric approximation can speed up the convergence but also introduces errors. Data fusion such as averaging will reduce communication costs but will also slow down the convergence (primarily because of repeated use of information in data fusion). In practice, we have to contend with a compromise between fast convergence, accurate information sharing and low storage and communication cost. Inspired by these findings, a problem-oriented hybrid protocol that takes the advantages of different approaches while minimizing the side-effects will be valuable.

VI. CONCLUSION

Flooding is an efficient albeit simple solution for information sharing over networks and is the basis of many other networking protocols. In this paper, we formulated it from a set-theoretic perspective, named distributed set-theoretic information flooding (DSIF). This led to a novel consensus protocol for networking referred to as collecting consensus, which has significant both advantages and disadvantages over averaging consensus and diffusion. We have analyzed the explicit convergence and optimality of DSIF based on a novel metric of DoC (degree of consensus). Practical solutions have been proposed either to determine the minimum number of iterations required for any desired DoC or to calculate the DoC that can be achieved by an actual number of iterations. It has also been noted that to save communication, data fusion (such as averaging) can be employed during flooding, which however may cause repeated information use and slower convergence. This trade-off has been analyzed.

Based on the theoretical results, a distributed particle filter framework is proposed and implemented for nonlinear target tracking which applies DSIF on sensor data alone or jointly with intermediate estimates. Simulations have demonstrated the convergence of the DSIF (faster than averaging consensus), the relationship between the filter performance and the DoC, and the advantage and disadvantage of applying parameterized approximation and data fusion for networking.

ACKNOWLEDGMENT

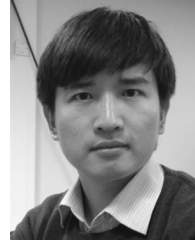
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as the likelihood consensus code from Dr. O. Hlinka. The authors would also like to acknowledge one anonymous reviewer's idea for a protocol which 1) applies DSIF for a certain number of time intervals, then 2) reduces the amount of stored information through a weighted averaging — like step, and 3) restarts with DSIF approach. We believe that this is a valuable research topic.

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Convergence of Distributed Flooding and Its Application for Distributed Bayesian Filtering

Tiancheng Li, Juan M. Corchado, *Member, IEEE*, and Javier Prieto, *Member, IEEE*

Abstract—Distributed flooding is a fundamental information sharing method to get network consensus via peer-to-peer communication. However, a unified consensus-oriented formulation of the algorithm and its convergence performance are not yet explicitly available in the literature. To fill this void in this paper, set-theoretic flooding rules are defined by encapsulating the information of interest in finite sets (one set per node), namely distributed set-theoretic information flooding (DSIF). This leads to a new type of consensus referred to as “collecting consensus,” which aims to ensure that all nodes get the same information. Convergence and optimality analyses are provided based on a consistent measure of the degree of consensus of the network. Compared with the prevailing averaging consensus, the proposed DSIF protocol benefits from avoiding repeated use of any information and offering the highest converging efficiency for network consensus while being exposed to increasing node-storage requirements against communication iterations and higher communication load. The protocol has been advocated for distributed nonlinear Bayesian filtering, where each node operates a separate particle filter, and the collecting consensus is pursued on the sensor data alone or jointly with intermediate local estimates. Simulations are provided in detail to demonstrate the theoretical findings.

Index Terms—Consensus, diffusion, distributed tracking, particle filter, sensor network.

I. INTRODUCTION

DISTRIBUTED computation has gained immense attention in the past decade, accompanying the rapid development and popularity of wireless sensor networks. In the successful networking operation, it is often of high interest that each node iteratively shares information with its intermediate neighbors (namely peer-to-peer communication) and consequently the entire network tends to reach a global alignment [1]/consensus [2]–[4] (to a certain degree). Compared to the centralized networking solutions based on a fusion center, distributed networking offers several advantages regarding scalability to adding or

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removing nodes, immunity to node failure, and dynamic adaptability to network topology changes.

However, there is a significant conflict between the degree of consensus (DoC) and communication requirement, as a higher DoC requires more communication, either more communicating iterations or higher communicating bandwidth which are limited by real time implementation and the communicating affordability of the nodes, respectively. Therefore, it is of paramount significance to seek a good balance for the trade-off so that the network achieves a satisfactory consensus in real-time and with affordable communication costs, which forms the majority of the research in the literature.

One of the most fundamental solutions for network information sharing is the flooding carried out in a distributed manner, by which all nodes synchronously broadcast their information to neighbors from the near to the distant. This protocol is well known in a few areas such as the communications [6]. Given that the network is strongly connected, it is able to achieve complete consensus (CC, i.e., all nodes have exactly the same set of information) after a certain number of iterations of peer-to-peer communication. This is quite appealing in theory but poses crucial challenges to the storage and communication requirement for large networks in practice. In fact, flooding has rarely been investigated in literature for distributed filtering (except for a few works, e.g., [7], [8]) even for small networks, regardless of its fast convergence (to be explicitly demonstrated in this paper) and ease of implementation. To note, there is one work that also refers to distributed flooding [9] which transmits the information of one single node over the network for routing. It is very different from the distributed protocol we consider here.

To date, a unified consensus-oriented formulation of the flooding algorithm and its convergence analysis are still missing in the literature. On one hand, there are cases in which DoC is required in the first priority while the sensor nodes have sufficient node-storage and communicating affordability to do so, for which flooding or even CC is simply preferable. On the other hand, instead of CC, it is more desired to perform flooding in a fewer, affordable, number of iterations for real time realization. What then can be expected and how can the number of iterations be properly determined?

To address the void and to answer the questions above, this paper aims to contribute from three aspects:

- 1) Formulate the flooding algorithm by encapsulating the information of interest in a finite set at each node and define set-theoretic flooding rules, namely *distributed set-theoretic information flooding* (DSIF), raising a new type of consensus termed *collecting consensus*.
- 2) Analysis the convergence and optimality of the DSIF protocol, based on a novel, consistent, metric of the DoC.

It is shown that DSIF enjoys the highest efficiency for network consensus among all distributed peer-to-peer communication schemes while suffering from heavier storage and communicational costs.

- 3) Show how the proposed DSIF scheme can be applied for and can benefit distributed Bayesian filtering, in which each node runs a separate particle filter (PF). Local PFs share sensor data alone or jointly with intermediate posteriors via DSIF. The latter usually needs to be parameterized in order to reduce the communication cost. The gain and loss to do so are analyzed and demonstrated in simulations.

The remainder of this paper is organized as follows. Notations and definitions regarding networking and three prevailing distributed information sharing protocols are given in Section II. As the main theoretical contribution, the collecting consensus-oriented DSIF is formulated in Section III with an analysis of its convergence and optimality. Section IV shows how to apply DSIF for distributed PF (DPF), with a brief literature review for DPF also given. Simulations are given in Section V and we conclude in Section VI.

II. NOTATION, DEFINITION AND BACKGROUND

A. Notation

The network topology is represented by a directed graph $G = (V, E)$ with the set of nodes $V = \{1, 2, \dots, N\}$ and the set of edges $E \subseteq V \times V$. In the directed graph, any edge is denoted by an ordered pair of nodes $(i, j) \in E$, which means node j is directly reachable from node i , where i is called the in-neighbor of j while j is the out-neighbor of i . For any $j \in V$, denote $N_j := \{i \in V | (i, j) \in E, i \neq j\}$, which is the set of all the in-neighbors of node j excluding node j itself. Undirected graph is a special type of directed graph where for any $(i, j) \in E$, we must have $(j, i) \in E$. If there exists a sequence of connected edges as follows

$$\{(i, \cdot), \dots, (\cdot, j)\} \subseteq E \quad (1)$$

This sequence of edges is called a path from node i to node j (denoted as $Path_{i \dots j}$) and node j is said to be “reachable” from node i . A digraph is said to be “strongly connected” (SC) if any node is reachable from all the other nodes, which is the digraph that can reach CC. For undirected graphs, SC is the same as connectivity [3].

The length of a path is given by the number of edges on that path. The length of the shortest path (perhaps, not unique) from node i to node j is called the distance from node i to node j , denoted as $D(i - j)$. Particularly, $D(i - i) = 0$ and if $(i, j) \in E$, $D(i - j) = 1$. We denote the set of all the nodes that are of distance $t \in \mathbb{N} = \{0, 1, 2, \dots\}$ to j as $N_j(t)$ and that of distance $t \in \mathbb{N}$ or smaller to j as $N_j(\leq t)$, namely $N_j(t) := \{n \in V | D(n - j) = t\}$, $N_j(\leq t) := \{n \in V | D(n - j) \leq t\}$. Obviously, we have $N_j(1) = N_j$, $N_j(0) = j$.

The largest distance between any two nodes, denoted as D_m , is called the diameter of the graph which is given by

$$D_m = \max_{i, j \in V} D(i - j) \quad (2)$$

Clearly, D_m only exist in SC networks. For any SC networks of at least two nodes, we have $D_m \in [1, N - 1]$, where the left bound corresponds to the fully connected network in which all nodes are in-neighbors of the others, while the right bound corresponds to the weakest connected network where all nodes are on a single chain in order, each having no more than two intermediate neighbors.

With particular regard to the distributed filtering problem, we assume that each node has independent abilities for: (i) filtering calculation, (ii) sensing to collect observations and (iii) communicating to neighbors. The communication is carried out in recursive iterations between neighboring nodes, each iteration consisting of sending no more than one data packet and receiving no more than one data packet. In addition, we need to clarify the following two definitions.

Definition 1 (real time communication): Communication is fully carried out between two successive observations and causes no sensor data missing or time-delay to the filter.

Definition 2 (communication bandwidth): The maximum size of the data packet that one node can send to or receive from its neighbor per communication.

B. Averaging/Maximum/Minimum Consensus

Here we assume that each node has a local scalar value, referred to as *state*, and it is of interest to compute the average of these values. An averaging consensus algorithm [2], [3] of zero communication time-delay is to reach an agreement regarding the state x_i each node has with local adapting dynamics $u_i(t)$, which can be written in discrete-time as

$$x_i(t + 1) = x_i(t) + u_i(t) \quad (3)$$

where $t \in \mathbb{N}$ denotes the communication iteration, $x_i(0)$ and $x_i(t)$ denote the initial and updated state of node i after iteration t , respectively.

In most of the time, the dynamics $u_i(t)$ is defined as

$$u_i(t) = \sum_{j \in N_i} \omega_{j \rightarrow i} (x_j(t) - x_i(t)) \quad (4)$$

where $\omega_{j \rightarrow i}$ is neighboring weight from node j to node i [2]–[4].

The complete convergence of *averaging consensus*: at iteration t states that, for any $i, j \in V$,

$$x_i(t) = x_j(t) \quad (5)$$

and asymptotically convergence (in the sense that $t \rightarrow \infty$)

$$\|x_i(t) - x_j(t)\| \leq \epsilon \quad (6)$$

where $\|x - y\|$ is a measure of the discrepancy between x and y , and ϵ is an error bound or margin [10].

In contrast to the average, one might be only interested in the maximum/minimum state, namely maximum/minimum consensus, which defines the iteration as

$$x_i(t + 1) = \max / \min \{x_i(t), \{x_j(t)\}_{j \in N_i}\} \quad (7)$$

The averaging consensus was well investigated in the community of control and systems [2]–[5]. One major concern is with the convergence, for which the spectral properties of the

183 graph Laplacian play a crucial role [2], [3], [10]-[12]. The sec-
 184 ond concern is raised by the information correlation/dependence
 185 among neighbors (especially when they own in part the same
 186 information). To account for this, fusion weights will be as-
 187 signed to coordinate the nodes, e.g., covariance intersection
 188 [14],[15], which has inspired many strategies. In the absence of
 189 clear information about the correlation or relative quality of the
 190 information among nodes, two weighting methods have been
 191 proposed: one is to weight all nodes equally [16] and the other
 192 is to weight them according to the size of their neighborhood
 193 (named Metropolis weights [17]).

194 Particularly for the tracking problem, the sensors may updates
 195 their observation frequently, preventing sufficient peer-to-peer
 196 communicating to get the network converge. This necessitates
 197 limiting the number of communicating iterations to gain a trade-
 198 off or compromise between a high DoC and little missing or
 199 time-delay of sensor data.

200 C. Gossip and Diffusion

201 To save communication, one alternative is to apply gossip
 202 to randomly choose fewer neighbors at each time (rather than
 203 to all neighbors) for averaging. It turns out that under mild
 204 conditions this process converges over time asymptotically [18].
 205 Gossip based distributed filtering has been reported in, e.g.,
 206 [19],[20]. However, gossip experiences the same problem as
 207 inefficient/repeated computations, for example, the same set (or
 208 largely similar set) of nodes repeatedly fusing their information
 209 at different points in time.

210 As another alternative, diffusion [22]-[24] performs only one
 211 iteration of peer-to-peer communication (i.e., the sensing and
 212 consensus time scales are the same), avoiding the problem of
 213 repeated use of any information. Based on it, the distributed
 214 Kalman filter (DKF) [23],[24] does not only share sensor data,
 215 but also local intermediate estimates through a diffusion update
 216 step. By this, the one-iteration-only communication is actually
 217 carried out on two types of data: the sensor data for the in-
 218 cremental update, and the estimates for the diffusion update.
 219 Hybrid fusion has been previously studied in the network using
 220 a fusion center, e.g., [21].

221 III. COLLECTING CONSENSUS AND DSIF

222 A. Collecting Consensus

223 In this paper, we are interested in an information sharing pro-
 224 tocol that does not repeatedly use any information, and that will
 225 converge to CC in a definite number of iterations. By CC, we
 226 mean that all the nodes have exactly the same information of
 227 interest. Such a consensus model in which each node aims to
 228 collect information from all reachable nodes via the shortest
 229 paths is referred to as collecting consensus. Different to aver-
 230 aging consensus, the information from different nodes remains
 231 conditionally independent (or more precisely stated, unfused)
 232 until the end of the communication, which requires that nodes
 233 have sufficient storage allowance.

234 To perform collecting consensus, the information that needs
 235 to be communicated is encapsulated as a set, and the DSIF
 236 algorithm defines the information set dynamics (in contrast to

(3)) based on the union operation

$$I_i(t+1) = I_i(t) \cup u_i(t) \quad (8)$$

238 where $I_i(t)$ and $u_i(t)$ denote the existing and new incom-
 239 ing information set of node i at iteration $t \in \mathbb{N}$ respectively,
 240 $I_i(0)$ denotes the initial information set at node i with the size
 241 $|I_i(0)| = 1, \forall i \in V$, and $u_i(0)$ the initial dynamics.

242 As a result of CC, all nodes shall have exactly the same
 243 information set, i.e., $\forall i \in V, t \geq D_m$,

$$I_i(t) = \bigcup_{j \in V} I_j(0) \quad (9)$$

244 To this end, the DSIF algorithm consists of two stages:

- 245 1) In the starting iteration, each node collects information
 246 from all its in-neighbors

$$u_i(0) = \bigcup_{j \in N_i} I_j(0) \quad (10)$$

- 247 2) In the following iterations $t \in \mathbb{N}^+ = \{1, 2, \dots\}$, each
 248 node collects the new information that its in-neighbors
 249 have received at the preceding iteration

$$u_i(t) = \bigcup_{j \in N_i} \{I_j(t) \setminus I_j(t-1)\} \quad (11)$$

250 where $A \setminus B$ is the set difference of A and B , namely the
 251 set of all elements that are members of A but not of B and
 252 when $t \geq D_m$, we will actually have $u_i(t) = \emptyset$.

253 As shown in (11), the receiving neighbors will sort out the new
 254 received data, which they then transmit to their out-neighbors
 255 in the next iteration. In this process, the same information may
 256 be repeatedly received over edges, leading to information over-
 257 use and communication power waste, which is one defect of
 258 the naive flooding protocol, named implosion [6]. To avoid this,
 259 we define the set-theoretic information flooding rules in the
 260 following.

261 B. Set-Theoretic Flooding Rules

262 *Rule 1 (data sending):* Each node only sends to its out-
 263 neighbors the new information that has never been flooded be-
 264 fore, and does so no more than once in each iteration.

265 *Rule 2 (data accepting):* Each node will not repeatedly take
 266 in the same information either from different in-neighbors or
 267 from the same node at different iterations, but only accept the
 268 information at its first arrival.

269 For both rules above, the data from each node shall be associ-
 270 ated with a unique ID for distinguishing. Given these two rules
 271 respected, we will have $\forall i \in V, t \in \mathbb{N}$,

$$|I_i(t)| = |N_i(\leq t)| \quad (12)$$

272 To combat time-increasing storage requirement and commu-
 273 nication load (when $t < D_m$), each element of data (often called
 274 a tuple) may be somehow compressed via e.g., dimension re-
 275 duction [25] and polynomial encoding [13], under the premise
 276 that little or even no information would be lost and the data from
 277 different nodes remain conditionally independent.

278 It is worth noting that in the case of maximum or minimum
 279 consensus, there is neither a problem of information set-size

growing nor information overuse as the fusion result is always a single maximum or minimum value.

C. Convergence and Optimality of DSIF

To gain insights of the convergence of the proposed DSIF scheme, we need a metric to measure the DoC for collecting consensus, for which we propose a metric based on the size of the information set as follows.

Definition 3 (DoC): The DoC, denoted as C^o , of a network with N nodes, is defined as follows

$$C^o(t) = \frac{\sum_{i=1}^N |I_i(t)| - N}{N(N-1)} \quad (13)$$

where $t \in \mathbb{N}$ denotes the number of DSIF iterations that has been performed and in the following we limit it to $t \leq D_m$. On the DoC, we have the following theorem, which states the convergence property of the DSIF protocol.

Theorem 1: $0 \leq C^o(t_1) < C^o(t_2) \leq 1, \forall 0 \leq t_1 < t_2 \leq D_m$.

Proof: Before performing DSIF, each node has its original one unit of data, i.e., $|I_i(0)| = 1, \forall i \in V$. That gives $C^o(0) = 0$. After $t \geq D_m$ DSIF iterations, CC will be reached as all nodes will have the same information, i.e., $|I_i(t)| = N, \forall i \in V$. Furthermore, we have the following two straightforward Claims (for which we omit any proof):

Claim 1: As stated by (12), the size of the information set owned by sensor $i \in V$ will not be reduced during flooding except that data fusion or removal is taken, i.e., $|I_i(t_1)| \leq |I_i(t_2)|$ for any $0 \leq t_1 < t_2$.

Claim 2: Supposing two nodes g and q are of distance D_m (namely $D(g-q) = D_m$), for any $0 < t \leq D_m$, there must exist at least one node $j \in N_q(t)$ on $Path_{g \dots j \dots q}$ satisfying $D(j-q) = t$ whose information will arrive to node q exactly at iteration t and then, we have $|I_q(t-1)| + 1 \leq |I_q(t)|$.

From these two claims, we may conclude that $C^o(t_1) < C^o(t_2), \forall 0 \leq t_1 < t_2 \leq D_m$, to accomplish the proof. ■

Theorem 2: In the sense of DoC as given in (13), the proposed DSIF achieves the highest converging efficiency among all distributed peer-to-peer communication schemes.

Proof: From the definition of the distance between nodes and the DSIF peer-to-peer communication rules, we have two additional straightforward Claims:

Claim 3: All nodes whose information can reach node i in t iterations of peer-to-peer communication belong to $N_i(\leq t)$.

Claim 4: All nodes $q \in N_i(\leq t)$ will surely flood their information to node i in t DSIF iterations.

A combination of Claims 3 and 4 indicates that the DSIF will gain the largest possible $|I_i(t)|$ for any $i \in V$ and $t \in \mathbb{N}$ as claimed, which entails the converging optimality. ■

D. Trade-off between DoC and Number of Iterations

For a given network topology, the DoC is uniquely determined by the number of DSIF iterations. In turn, one can also determine the required number of iterations for a desired DoC, e.g., $T_c = 0.5$. That is, the DSIF stops at iteration t once

$$C^o(t) \geq T_c \quad (14)$$

Algorithm 1: DSIF operations at node i for DoC T_c .

INITIALIZATION:

1: $t \leftarrow 1; C^o(0) \leftarrow 0$

2: $I_i(t) = I_i(t-1) \cup \bigcup_{j \in N_i} I_j(t-1)$

RECURSIVE FLOODING ITERATION:

3: While $C^o(t) < T_c$

4: $t \leftarrow t + 1$

5: $I_i(t) \leftarrow I_i(t-1) \cup \bigcup_{j \in N_i} \{I_j(t-1) \setminus I_j(t-2)\}$

6: $C^o(t) \leftarrow \frac{\sum_{i=1}^N |I_i(t)| - N}{N(N-1)}$

7: End while

8: Return: $I_i(t)$

Algorithm 1 summarizes the communicating operations that need to be performed on node i for a given DoC T_c . 329 330

For a constant network of a known topology, the minimum number of iterations can be determined a priori by (14). However for time-varying dynamic networks, it needs to be calculated online via a consensus algorithm. To facilitate the use in time-varying networks without burdening any consensus procedures, we define the local DoC metric as follows: 331 332 333 334 335 336

Definition 4 (Local DoC): The DoC of node $i \in V$, denoted as C_i^o , after t DSIF iterations, is given as 337 338

$$C_i^o(t) = \frac{|I_i(t)| - 1}{N_k - 1} \quad (15)$$

where N_k is number of nodes in the network at time k , which needs to be estimated if unknown. From here we derive the following theorem as the local-node version of Theorem 1. 339 340 341

Theorem 3: $0 \leq C_i^o(t_1) \leq C_i^o(t_2) \leq 1, \forall 0 \leq t_1 < t_2 \leq D_m$ 342

Proof: This theorem states the same content as Claim 1. As a key difference to Theorem 1, the equality $C_i^o(t_1) = C_i^o(t_2)$ holds when and only when the number of nodes that are of distance t_1 to node $i \in V$ is the same as that of distance t_2 , i.e., $|N_i(t_1)| = |N_i(t_2)|$. ■ 343 344 345 346 347

Setting a threshold, e.g., $T_c = 0.5$ which can be the same or different for different nodes, on the desired local DoC, the consensus updating at node i may stop at iteration t once 348 349 350

$$C_i^o(t) \geq T_c \quad (16)$$

Furthermore, based on DoC we can define the convergence speed (CoS), either globally or locally, to measure the change of the size of the information set at each iteration, which also indicates the local real-time communication bandwidth. 351 352 353 354

Definition 5 (CoS): At iteration $t \in \mathbb{N}$, the global CoS, denoted as C^s , is defined as, 355 356

$$C^s(t) = C^o(t) - C^o(t-1) \quad (17)$$

and the local CoS of node $i \in V$ is defined as 357

$$C_i^s(t) = C_i^o(t) - C_i^o(t-1) \quad (18)$$

Theorem 4: $C^s(t) > 0, C_i^s(t) \geq 0, \forall 1 \leq t \leq D_m$ 358

Proof: The theorem is immediate from Theorems 1 and 3 as for any $1 \leq t \leq D_m, i \in V$, we have $C^s(t) = C^o(t) - C^o(t-1) > 0$ from Theorem 1, and $C_i^s(t) = C_i^o(t) - C_i^o(t-1) \geq 0$ from Theorem 3. ■ 359 360 361 362

Theorems 1, 3 and 4 entail an appealing property of the DSIF protocol which will not only converge definitively, but also 363 364

365 has a guaranteed converging speed that is globally positive and
 366 locally non-negative everywhere and at any iteration until CC
 367 is reached. We refer to this as *strong convergence*. It, however,
 368 also indicates a (non-negative) increasing storage requirement
 369 against communicating iterations. As an alternative to (16), we
 370 can build the predetermined threshold on the local CoS, e.g.,
 371 $T_s = 0.1$, then the minimum number of iterations t needs to
 372 satisfy

$$C_i^s(t) \leq T_s \quad (19)$$

373 But it is critical to note that we do not have any monotonicity
 374 on the CoS, e.g., $C^s(t_2) \leq C^s(t_1)$ or $C_i^s(t_2) \leq C_i^s(t_1)$ for $1 \leq$
 375 $t_1 < t_2 \leq D_m$. Therefore, the CoS at iteration t does not say
 376 anything of the CoS at iteration $t + 1$.

377 E. Comparison and Practical Consideration

378 Both metrics of DoC and CoS are clearly defined and easier
 379 to calculate than the one proposed for averaging consensus, e.g.,
 380 convergence rate [10]–[12], steady-state mean-square deviation
 381 [4] or disagreement vector [2], [3]. As indicated by Theorem 2,
 382 no peer-to-peer communication protocols converge faster than
 383 DSIF in terms of DoC. This superiority, however, is achieved
 384 at the expense of higher node storage requirements and heavier
 385 communication bandwidths. If the size of the data set at one
 386 node exceeds its communication bandwidth, multiple iterations
 387 will then be needed for that data set, otherwise data fusion is
 388 required to control the data size. In the former case, the required
 389 number of iterations will increase, while in the latter case the
 390 information completeness or independence may not be kept.
 391 However, we will not address this issue further here, which is
 392 quite problem dependent. In brief, we have the following remark
 393 on the respective advantages of averaging consensus, diffusion
 394 and collecting consensus.

395 *Remark 1:* The averaging consensus takes the lowest com-
 396 municating bandwidth (always one unit of data) but more iter-
 397 ations to reach any DoCs while the diffusion severely limits
 398 the number of iterations (to one only) which may insufficiently
 399 use the communication affordability (i.e., more iterations are
 400 actually allowed in real time communication). In contrast, the
 401 proposed DSIF protocol aims to get the best possible consensus
 402 in an real-time-allowed number of iterations, which is therefore
 403 particularly suited to small and moderate networks for which
 404 the nodes have sufficient storage and communicating power. A
 405 means to facilitate its use in large networks is to selectively
 406 apply data fusion such as averaging in every several flooding
 407 iterations in order to control the data-set size. This will lead to
 408 a hybrid protocol that iterates between flooding and averaging
 409 consensus, to gain a balance between benefiting from high com-
 410 munication efficiency and suffering from information overuse
 411 and slower convergence.

412 IV. DISTRIBUTED BAYESIAN FILTERING USING DSIF

413 A. State-of-the-art DPF Protocols

414 Before presenting our DPF framework based on DSIF, a brief
 415 revisit of the PF algorithm and existing DPF protocols is given

below. Suppose that at time k , the local (marginal) posterior at
 sensor i is represented by a local PF

$$p(\mathbf{x}_k | \mathbf{z}_{i,1:k}) \approx \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} \delta(\mathbf{x}_k - \mathbf{x}_{i,k}^{(m)}) \quad (20)$$

where $\delta(\mathbf{x} - \mathbf{y})$ is the Dirac delta impulse, which equals to
 one if $\mathbf{x} = \mathbf{y}$ and to zero otherwise, \mathbf{x}_k is the true state vector,
 $\mathbf{z}_{i,1:k}$ is the observation serial, $\mathbf{x}_{i,k}^{(m)}$ and $w_{i,k}^{(m)}$ are the state and
 normalized weight of the m th particle respectively, $M_{i,k}$ is the
 total number of particles at filtering time k .

The essence of the PF is to assess how well each particle
 conforms to the state model and explains the observations, using
 this assessment to generate a weighted sample approximation
 to the Bayesian posterior, and thereby form sub-optimal state
 estimates. Given local measurement $\mathbf{z}_{i,k}$, $i \in V$, the weights of
 the particles are evaluated over time based on the sequential
 importance sampling (SIS) principle as

$$w_{i,k}^{(m)} \propto w_{i,k-1}^{(m)} \frac{p(\mathbf{z}_{i,k} | \mathbf{x}_{i,k}^{(m)}) p(\mathbf{x}_{i,k}^{(m)} | \mathbf{x}_{i,k-1}^{(m)})}{\pi(\mathbf{x}_{i,k}^{(m)} | \mathbf{x}_{i,k-1}^{(m)}, \mathbf{z}_{i,1:k})} \quad (21)$$

where $\pi(\cdot)$ is a proposal to generate particles, and in general
 its design shall take into account both the newest measure-
 ment $z_{i,k}$ and the prior in order to best match the posterior; see
 e.g. [19], [28], [31]. The use of the observation in the sampling
 proposal design is particularly helpful (and even necessary for
 avoiding sample degeneracy) when the observation is very ac-
 curate. However, caution should be exercised here since the
 repeated use of the observation (both for proposal design and
 in likelihood calculation) may not benefit the filter when the
 observation suffers from significant noise [42].

In addition to SIS, resampling is usually required to reduce
 the weight variance when it exceeds a certain threshold, so that
 all particles will have equal or approximate weights while the
 posterior distribution can be the best maintained [31], [32]. This
 has often been referred to as sampling importance resampling
 (SIR), which is the core of the majority of existing PFs. We
 assume the reader is familiar with the centralized PF and so
 limit ourselves hereafter to the distributed implementation, in
 which local nodes carry out PF calculations in parallel and
 meanwhile share information with their neighbors to assist their
 filters. For this, a variety of information sharing protocols have
 been proposed, which can be classified as follows:

- 1) *Sequential information passing:* Information transmits in
 a sequential, predefined manner from a node to one of
 its neighboring nodes via a cyclic path until the entire
 network is traversed [43]. The sequential realm is sensitive
 to the mobility and failure of nodes/edges and is time-
 consuming.
- 2) *Flooding:* As addressed, the flooding protocol provides
 the fastest albeit communication-intensive way to spread
 information over the network [7], [8], but, neither any clue
 to determine the number of communication iterations in
 order to compromise real time realization and DoC nor
 any convergence results has been shown.

3) *Averaging consensus*. There is a large body of work concerning averaging consensus-based distributed filtering. The data transmitted between neighboring nodes can be posterior statistics in the form of Gaussian component [33] /GM [29]–[30] or generalized probability densities [36]–[37], likelihood [26]–[28], particle set [34]–[35] or raw observations [38]. Excellent surveys are also available such as a taxonomy of DPFs [39], a comparison of several belief consensus algorithms [40] and a recent survey of convergence and error propagation of DPFs [28]. In summary, complete information sharing affords better accuracy but has higher communication requirements, such as [34]–[35] that exchange all particles. Parameter approximation [26]–[33] or random gossip [19]–[20] can significantly reduce the communication cost, but may lead to a deterioration in the filter performance.

4) *Diffusion*: The diffusion scheme addressed in Section II.C also provides a competitive alternative to the averaging consensus for DPF [7]–[8], [41].

We note that the sensor data can be either simple (e.g., range, bearing) or complex (e.g., image data). To avoid distracting from the key contribution of this paper on collecting consensus and the DSIF protocol, we only consider the former case for simplicity. For the latter case, one may consider compressing the sensor data, e.g., [25]–[26], [12]–[13] or transmitting the low-dimensional likelihood for replacement [26]–[27]. At the current stage, we have not considered complicated network issues such as communication constraints, e.g., [45]–[46], and asynchronous sensing, e.g., [47]–[48]. However, we note all of these issues are valuable to be investigated on the base of the proposed DSIF protocol.

B. DSIF on Sensor Data and on Local Posterior

In the proposed DPF framework, the DSIF scheme will be applied on the sensor data alone or jointly on local posteriors. In the latter, we propose parameterizing the posterior to save communication. Since a vast number of random numbers are required by the PF, it is communication intensive to run consensus on them, and it is not our intention to do so.

First, DSIF is implemented on the sensor data including the target-observations (and uncertainties) associated with the sensor ID, all as one unit. To note, the sensor position is often required for likelihood calculation and therefore can serve as the unique sensor ID for distinguishing. Then, the resultant consensus on sensor data with sensor profiles given a priori, is equivalent to collecting consensus on the likelihood which is required for PF updating. A likelihood function contains the information of both the sensor data and the sensor profile in a more compact manner. But for simplicity of understanding, we keep addressing consensus on sensor data.

The filtering posteriors obtained at different nodes, referred to as local posteriors, will be different, even if CC is reached on sensor data over the network where the difference attributes to the different random numbers. If DoC is low on sensor data, the difference between local posteriors will be relatively significant. As such, we may apply the second DSIF scheme to fuse local posteriors among neighbors as well as to get the local LMS

(least mean squares) estimate; we refer to this step as diffusion, in parallel to [24]. By this, each node aims to improve their local estimate with regard to their neighbors' posterior. However, parameter approximation of local posteriors, typically via Gaussian or GM approximation, is needed (otherwise massive communication will be triggered if the complete posterior is communicated by transmitting the entire particle set), which will in turn introduce approximation errors to the posterior. This trade-off is much problem-dependent and will determine whether the second DSIF is worthwhile.

The operations that need to be conducted on each sensor in the proposed distributed PF is summarized in Algorithm 2. In it, steps 1-a and 1-b are independent of each other and therefore can be carried out in either order or in parallel. Sensor data DSIF and posterior DSIF have been implemented t_1 and t_2 iterations respectively, where t_1 and t_2 are not necessarily equal but are determined for respective desired or the largest affordable DoCs as addressed in Section III. They show complementary features and resemble the Incremental and Diffusion updates of the diffusion-based DKF [24]. But, there are obvious differences:

- 1) Our framework is developed for nonlinear models which releases the requirement of linear system functions and even Gaussian assumption of the posterior;
- 2) Our consensus protocol does not limit information sharing between neighbors to one iteration only but instead, the DoC will be pursued as much as the real time communication allows;
- 3) Our diffusion update (Step 5 in Algorithm 2) is an optional step, which is advocated for re-setting local posteriors only when local posteriors are significantly different (as a consequence of a low DoC on the sensor data achieved in the first DSIF implementation). When the difference between local posteriors is insignificant (because of a high DoC achieved on the sensor data), there will be less need to further fuse them and so it may be better not to diffuse local posteriors since the errors introduced due to parameterization can be more significant than the benefit. This is a critical point. We will demonstrate this in detail through simulations in Section V. In addition, we provide two easy-to-implement diffusion choices.
- 4) We point out that the proposed two DSIF procedures can be performed jointly, although this may not reduce the communication load and the storage requirement in total; see the following Remark 2.

Remark 2: Two DSIF implementations regarding the sensor data and the local filter estimates form the starting step and the end step of each filtering iteration, respectively. In the time series, they are adjacent. Therefore, they may be combined in one joint consensus scheme at some stages (which however does not necessarily indicate that $t_1 = t_2$), i.e., the local estimates obtained at filtering time k can be combined with sensor data received at time $k + 1$ as one unit of data, both sharing the same node ID for DSIF. Then, the initial information set at node $i \in V$ can be defined as

$$I_{i,k}(0) := \{\hat{\mathbf{x}}_{i,k-1}, P_{i,k-1}\} \cup \mathbf{z}_{i,k} \quad (22)$$

Algorithm 2: Distributed PF calculation executed on node i .

Step 1-a Filter prediction: Propagate the particles $\mathbf{x}_{i,k-1}^{(m)}$ to $\mathbf{x}_{i,k}^{(m)}$ for all $m = 1, \dots, M_{i,k}$, according to the proposal function $\pi(\mathbf{x}_{i,k}^{(m)} | \mathbf{x}_{i,k-1}^{(m)}, \mathbf{z}_{i,1:k})$. At $k = 0$, particles are sampled from an initial proposal π_0 instead for filter initialization.

Step 1-b 1st DSIF: Perform t_1 DSIF iterations on sensor data as given in Algorithm 1, resulting in a combined measurement set $\mathbf{Z}_{i,k} = \{\mathbf{z}_{j,k}\}_{j \in N_i(\leq t_1)}$. This step is carried out whenever new measurements become available.

Step 2 Filter updating: Re-weight all particles via (21) ($\mathbf{z}_{i,t}$ therein shall be replaced by $\mathbf{Z}_{i,k}$ obtained in the 1st DSIF) and then normalize them as follows

$$w_{i,k}^{(m)} \leftarrow w_{i,k}^{(m)} / \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} \quad (23)$$

Step 3 Estimate extraction: Extract local estimate $\hat{\mathbf{x}}_{i,k}$ and calculate their covariance $P_{i,k}$ from the local random measure $\chi_{i,k} = \{\mathbf{x}_{i,k}^{(m)}, w_{i,k}^{(m)}\}_{m=1,2,\dots,M_{i,k}}$ as follows

$$\hat{\mathbf{x}}_{i,k} = \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} \mathbf{x}_{i,k}^{(m)} \quad (24)$$

$$P_{i,k} = \sum_{m=1}^{M_{i,k}} w_{i,k}^{(m)} (\mathbf{x}_{i,k}^{(m)} - \hat{\mathbf{x}}_{i,k}^{(m)}) (\mathbf{x}_{i,k}^{(m)} - \hat{\mathbf{x}}_{i,k}^{(m)})^T \quad (25)$$

Step 4 2nd DSIF: Perform t_2 DSIF on local estimates obtained in Step 3, resulting in a set of intermediate estimates $\{\hat{\mathbf{x}}_{j,k}, P_{j,k}\}_{j \in N_i(\leq t_2)}$, which will be fused in the LMS sense as follows

$$\hat{\mathbf{x}}_{i,k}^{\text{LMS}} = \sum_{j \in N_i(\leq t_2)} \hat{\mathbf{x}}_{j,k} P_{j,k}^{-1} \quad (26)$$

$$P_{i,k}^{\text{LMS}} = \left(\sum_{j \in N_i(\leq t_2)} P_{j,k}^{-1} \right)^{-1} \quad (27)$$

This also offers the local filter output at sensor i .

Step 5 Diffusion: As an option, the shared filter estimates given in Step 4 can be used to re-set the local PF posterior $\chi_{i,k}$. If so, there are two choices (the second is expected to have a higher approximation accuracy than the first).

1) Re-set $\chi_{i,k}$ as the LMS fused Gaussian distribution, as is done in the Gaussian PF [49].

$$\chi_{i,k} \leftarrow \mathcal{N}(\hat{\mathbf{x}}_{i,k}^{\text{LMS}}, P_{i,k}^{\text{LMS}}) \quad (28)$$

2) Re-set $\chi_{i,k}$ as the shared GM before performing LMS fusion, as is done in the Gaussian sum PF [50]

$$\chi_{i,k} \leftarrow \sum_{j \in N_i(\leq t_2)} \mathcal{N}(\hat{\mathbf{x}}_{j,k}, P_{j,k}) \quad (29)$$

Step 6 Resampling: Sample from the updated particle set [32] if the variance of weights exceeds a specified threshold and if Step 5 is not applied. If Step 5 is applied, sample from the diffused Gaussian or GM distribution $\chi_{i,k}$ given by Step 5 to generate a new particle set. Update $k \leftarrow k - 1$ and go to the next filtering iteration.

V. SIMULATIONS

575

In this section, we consider tracking a target that moves in the $x - y$ plane by using the proposed DPF based on a constant sensor network earlier appeared in [39] as given in Fig. 1. The network has totally 10 sensors and a diameter $D_m = 4$. The simulation models and parameters are the same to [27]. In specific, we have the initial state as $\mathbf{x}_0 = [4, 0.5, 4, 0.5]^T$. The Markov transition model that governs the target movement of nearly constant velocity is given by

$$\mathbf{x}_k = \begin{bmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix} \mathbf{x}_{k-1} + \begin{bmatrix} 0.5 & 0 \\ 1 & 0 \\ 0 & 0.5 \\ 0 & 1 \end{bmatrix} \mathbf{u}_k \quad (30)$$

where $\mathbf{x}_k = [p_{x,k}, \dot{p}_{x,k}, p_{y,k}, \dot{p}_{y,k}]^T$, $[p_{x,k}, p_{y,k}]^T$ gives the position and $[\dot{p}_{x,k}, \dot{p}_{y,k}]^T$ the velocity, $\mathbf{u}_k \sim \mathcal{N}(\mathbf{0}_2, 0.00035\mathbf{I}_2)$.

The target emits an acoustic or radio signal with a known constant transmit power P_t that can be received by all sensors independently, i.e., the scalar measurement function of sensor i located at $[s_{i,x}, s_{i,y}]^T$ about target \mathbf{x}_k is

$$z_{i,k} = \frac{\alpha P_t}{\| [p_{x,k}, p_{y,k}]^T - [s_{i,x}, s_{i,y}]^T \|^{\gamma}} + v_k \quad (31)$$

where α is a constant that depends on several factors such as fast and slow fading, and gains in the transmitter and receiver antennas, γ is the path loss exponent [44], and $v_k \sim \mathcal{N}(0, \sigma_v^2)$ is the measurement noise. In parallel to [27], we set simply $\alpha P_t = 10, \gamma = 1, \sigma_v^2 = 0.001$.

When multiple synchronous observations are available, the weight of particles is updated by multiplying the likelihoods given by each available measurement. That is,

$$p(Z_{i,k} | \mathbf{x}_{i,k}^{(m)}) = \prod_{j \in N_i(\leq t)} p(z_{j,k} | \mathbf{x}_{i,k}^{(m)}) \quad (32)$$

where $Z_{i,k} = \{z_{i,k}\}_{j \in N_i(\leq t)}$ is the measurement set at sensor i gained in the first DSIF procedure of total t iterations.

For any sensor $i \in V$, the necessary and sufficient number of iterations, denoted as $D_{m,i}$, to receive the information from all the other sensors can be given by

$$D_{m,i} := \max_{j \in V} D(j - i) \quad (33)$$

We design three groups of simulations in the following three subsections that use the same ground truths to evaluate or compare the following five PF protocols, where the first three are distributed while the last two are centralized. All PFs use the same number of particles ($M = 1000$).

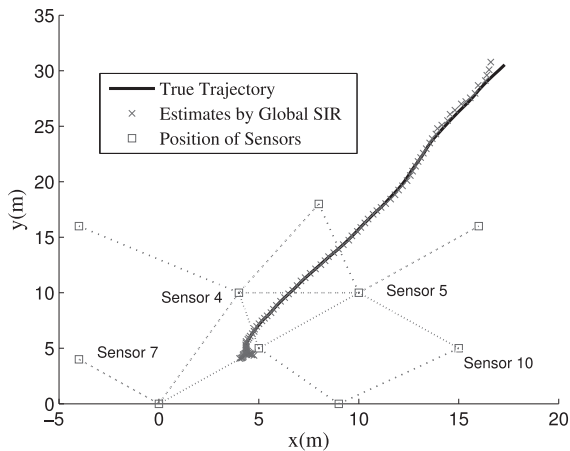


Fig. 1. The topology of the sensor network, the target trajectory and its estimate given by a global SIR filter in one trial.

- 1) *C-SIR*: we apply DSIF only on the sensor data, named Consensus without Diffusion (i.e., Steps 4 and 5 are not applied in Algorithm 2). In this case, each local PF is a SIR filter that is free of any Gaussian assumption;
- 2) *CD-GMPF*: we apply DSIF on both sensor data and local estimates named Consensus with Diffusion (i.e. Steps 4 and 5 are applied in Algorithm 2). In this case, each local PF is a Gaussian sum PF that applies (29) for posterior approximation and fusion;
- 3) *L-C-SIR*: the Likelihood Consensus-based SIR filter [27] can be viewed as a special case of our C-SIR filter that applies sensor data averaging consensus (for likelihood multiplying) at each iteration. For fast converging, the Metropolis weights strategy [52] is employed for averaging in the L-C-SIR filter;
- 4) *Local-SIR/GMPF*: local SIR filter or GMPF that does not communicate with each other at all;
- 5) *Global-SIR*: a centralized SIR filter that is able to access all sensor observations at all times.

To mitigate the problem of sample impoverishment that is often caused by resampling in the SIR filters, the minimum-sampling-variance resampling [32] is applied when and only when the effective sample size is smaller than $M/2$ and if applied, a roughening noise that is equivalent to half of \mathbf{u}_k will be used [31].

To measure the filtering accuracy, we calculate the root mean square error (RMSE) on both the position estimate and the velocity estimate, respectively, as follows

$$\text{RMSE}_{\text{pos}_k} = \sqrt{\frac{1}{C} \sum_{c=1}^C (x_{k,c} - \hat{x}_{k,c})^2 + (y_{k,c} - \hat{y}_{k,c})^2} \quad (34)$$

$$\text{RMSE}_{\text{vel}_k} = \sqrt{\frac{1}{C} \sum_{c=1}^C (\dot{x}_{k,c} - \hat{\dot{x}}_{k,c})^2 + (\dot{y}_{k,c} - \hat{\dot{y}}_{k,c})^2} \quad (35)$$

where $[\hat{x}_{k,c}, \hat{y}_{k,c}]^T$ and $[\hat{\dot{x}}_{k,c}, \hat{\dot{y}}_{k,c}]^T$ are the position-estimate and velocity-estimate given at filtering time k in trial c , respectively, and $C = 20$ is the total number of MC trials. Further, the

TABLE I
DoC ACHIEVED AT EACH DSIF ITERATION (LOCAL AND GLOBAL)

	Sensor 4	Sensor 5	Sensor 7	Sensor 10	Global
$t = 0$	0	0	0	0	0
$t = 1$	5/9	5/9	1/9	2/9	26/90
$t = 2$	1	8/9	3/9	6/9	61/90
$t = 3$	1	1	7/9	8/9	86/90
$t = 4$	1	1	1	1	1

average position RMSE is defined as the mean of $\text{RMSE}_{\text{pos}_k}$ over the entire simulation period of 100 filtering iterations. In each trial, the ground truth is independently generated (for generality). In all trials, the prior distribution of the particle set is initialized around the true state as $\mathcal{N}(x_0, P_0)$, with $P_0 = \text{diag}[2, 0.001, 2, 0.001]^T$.

In particular, we will assess the filter performance at four representative sensors, marked in Fig. 1 as sensors 4, 5, 7 and 10. For them, we have $D_{m,4} = 2, D_{m,5} = 3, D_{m,7} = 4, D_{m,10} = 4$. This means that sensor 4 will achieve CC first (after 2 iterations) while sensors 7 and 10 will be the last (after 4 iterations). For different numbers of DSIF iterations, the global and local DoCs are given in Table I. Particularly, for $t = 1$, we have the global DoC determined as

$$C^o(1) = \frac{|E|}{N(N-1)} \quad (36)$$

where $|E|$ is the number of edges; (a, b) and (b, a) are counted as two different edges.

A. Consensus without Diffusion

In this case, each sensor operates a separate SIR filter. Sensors are assumed conditionally independent and use different random numbers. The posteriors obtained by sensors will be different from both each other and the global/local PF, even given that they all reach CC on sensor data.

For different numbers of DSIF iterations from 0 (no consensus at all) to 4 (D_m), the RMSEs of the position and velocity estimation of local C-SIR filters and the global SIR filter are given in Fig. 2.(a)-(e) respectively, corresponding to different DoCs. The average RMSEs over 100 filtering steps against the number of DSIF iterations are given in Fig. 2.(f). The results clearly demonstrate that:

- 1) A single passive sensor is not capable of delivering good tracking in this problem as the RMSEs given by local PFs are much higher than that provided by DPFs; this necessitates the collaboration of multiple geographically dispersed sensors;
- 2) The more informative sensor data used, the better the filter performance;
- 3) The larger DoC, the closer the local PF performance to the centralized PF, i.e., local filters converge to the global filter against iterations as the DoC increases;
- 4) Once CC is reached, the performance of the local PF is very close to that of the centralized PF (with regard to

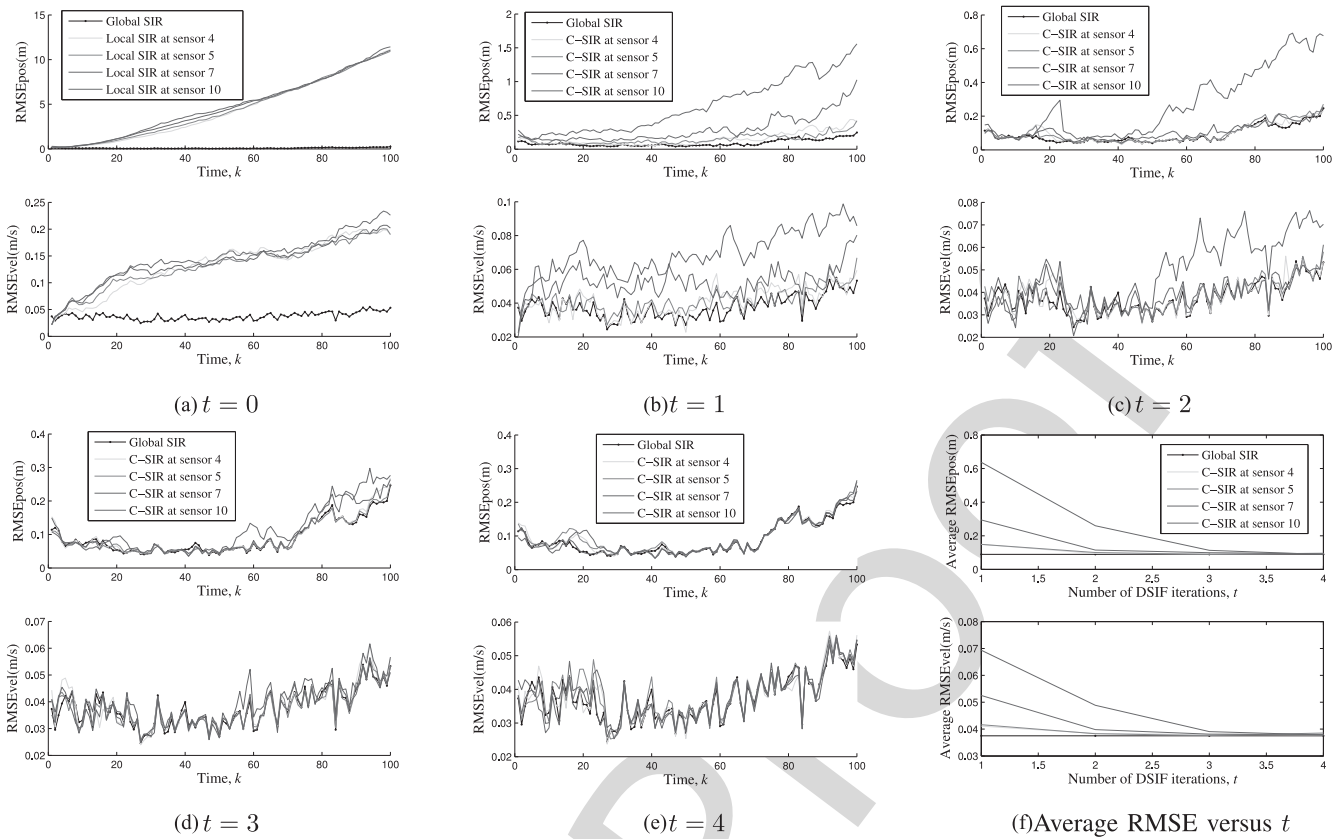


Fig. 2. Position and velocity RMSE of C-SIR filters with different numbers of DSIF iterations, comparing with the global SIR filter.

680 both position and velocity) but still not the same, since
 681 different random numbers are used.

682 Based on the measure of DoC, we are able to approximately
 683 determine how much information divergence different nodes
 684 will have and what payoff can be expected if one more or one
 685 less iteration of peer-to-peer communication is employed. For
 686 example, when the number of iterations is $t = 3$, the global
 687 DoC is as high as 86/90, close to 1, which agrees with the slight
 688 difference between Fig. 3(a) and (b). This is a valuable part of
 689 the metric of DoC.

690 B. Consensus with Diffusion

691 In this case, each sensor runs a separate GMPF. Collecting
 692 consensus are applied on both the sensor data and intermediate
 693 estimates jointly in a single DSIF procedure (and set $t_1 = t_2$).
 694 Because of the GM diffusion of intermediate estimates, the
 695 local fused estimates are expected to be closer to each other.
 696 If CC is reached, they shall be exactly the same. In parallel to
 697 the last simulation, different numbers of DSIF iterations from
 698 0 to 4 are employed to the CD-GMPFs, which are compared with
 699 the (centralized) global SIR PF in Fig. 3(a)–(e) respectively.
 700 The average RMSEs of these filters against the number of DSIF
 701 iterations are given in Fig. 3(f).

702 We use the same ground truth (20 MC trials) regarding the
 703 trajectories and sensor observations as the last simulation. Com-
 704 pared to the last simulation, we can find that

705 1) A single passive sensor can still hardly work well when
 706 the local SIR filters are replaced by local GMPFs;

- 707 2) Given the same number of DSIF iterations $t = 1, 2$, CD-
 708 GMPFs perform much better than C-SIR and are much
 709 closer to each other; this is because of the second DSIF
 710 scheme on the posteriors over the network which enhances
 711 the consensus to improve local estimates;
- 712 3) Given $t = 3, 4$ iterations, the local CD-GMPFs perform
 713 almost the same but different to the global SIR filter;
- 714 4) Given CC achieved, the RMSEs of all local GMPFs are
 715 exactly the same but are inferior to the global SIR fil-
 716 ter, especially at the later stage in this tracking example.
 717 Analysis and discussion will be given next.

718 C. Comparison and Discussion

719 Finally, we compare both types of DSIF-based DPFs with the
 720 L-C-SIR filter [27], [39]. The key difference of the likelihood
 721 consensus to DSIF is that each node fuses information interme-
 722 diately after receiving them and therefore the communication
 723 cost is lower, but it is exposed to repeated use of information
 724 and slower convergence.

725 First, for $t = 4$, the average (over all nodes) position RMSEs
 726 of the C-SIR, CD-GMPF and the L-C-SIR filters are given
 727 in Fig. 4. It shows that the C-SIR filter achieves the closest
 728 performance to that of the centralized filter. We further calculate
 729 the mean of these average RMSEs for $t = 0$ to 8 and the results
 730 are given in Fig. 5. It shows that these consensus protocols can
 731 all significantly improve the filter performance as compared to
 732 the local filter that applies no consensus and converges against
 733 communication iterations.

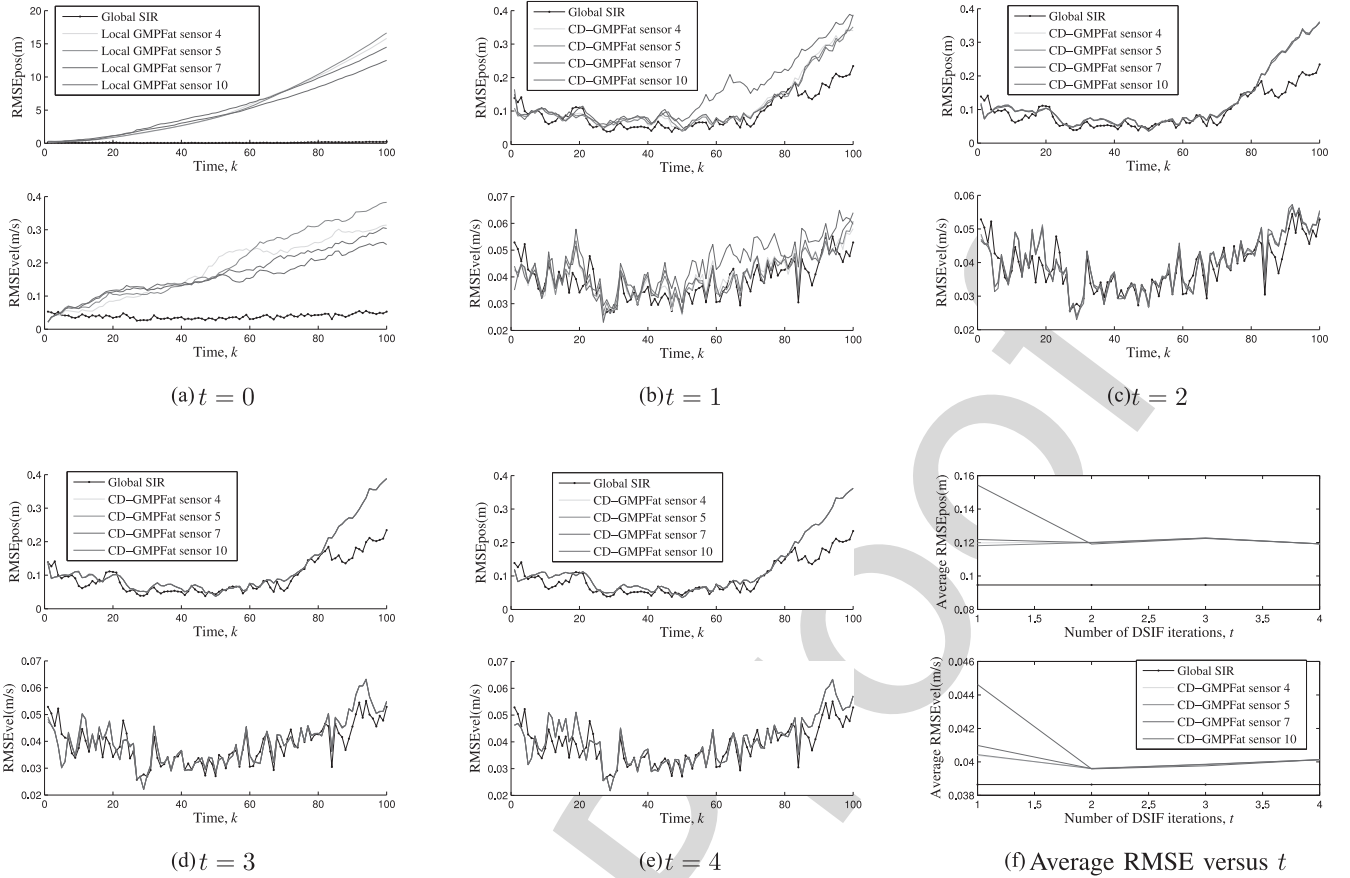


Fig. 3. Position and velocity RMSE of CD-GMPFs with different numbers of DSIF iterations, comparing to the global SIR filter.

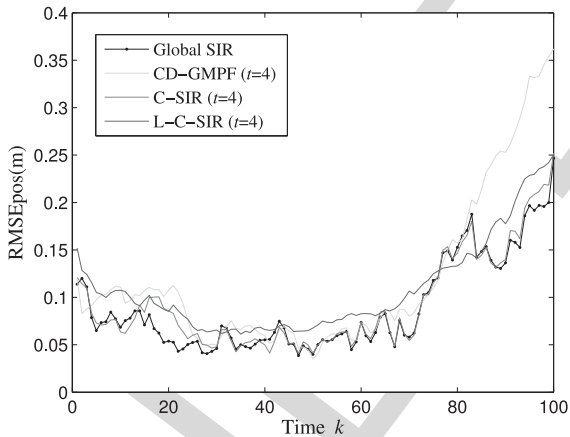


Fig. 4. Position RMSE of different DPFs applying 4 iterations of peer-to-peer communication.

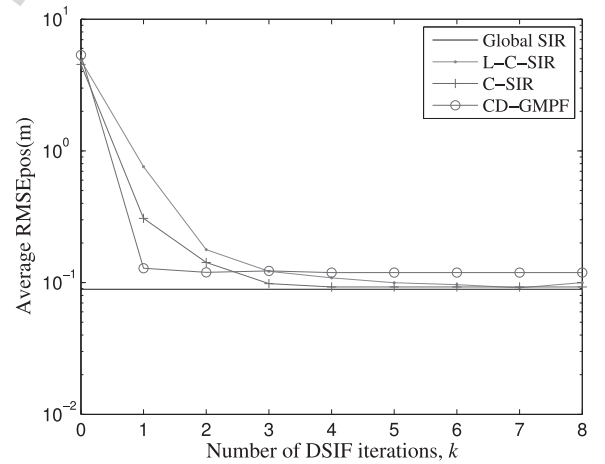


Fig. 5. Average position RMSE of different DPFs over 100 filtering steps against the number of peer-to-peer communication iterations.

734 Furthermore, we have the following observations, which
735 show more insights of these three types of DPFs:

- 736 1) DSIF based C-SIR and CD-GMPF converge faster than
737 the averaging consensus-based L-C-SIR filter at the ex-
738 pense of higher communication cost. CD-GMPF con-
739 verges the fastest but it suffers from a larger RMSE at
740 the end, all due to its diffusion step that shares infor-
741 mation among nodes more thoroughly than without diffu-
742 sion but also introduce errors;

- 2) For a relatively small number of iterations that correspond
743 to a low DoC on observation (which may lead to a large
744 discrepancy between local nodes' posteriors), the C-SIR
745 filter is inferior to the CD-GMPF, as shown in Fig. 5 (also
746 told by comparing between Figs. 2(f) and 3(f)). In this
747 case, the diffusion update leads to earlier convergence
748 and better performance for the filter. This is in line with
749 the findings reported in [23];
750

3) For a large number of iterations that correspond to a high DoC on observation and consequently on posterior (leaving little space to benefit from posterior fusion), the diffusion update of the CD-GMPF is not so preferable; instead, the GM approximation error caused in the diffusion might be more significant than the benefit it can offer, resulting in an overall filter degradation. We must note that if the whole particle sets are transmitted for diffusion without any approximations, and also the dependence between the posteriors are accounted for properly in the diffusion update, it shall always be beneficial in theory regardless of the much greater cost in communication and local fusion calculation.

These results confirm our theoretical prediction and demonstrate further that, both approximation and data fusion during communication can be either beneficial or counterproductive. Generally speaking, parametric approximation can speed up the convergence but also introduces errors. Data fusion such as averaging will reduce communication costs but will also slow down the convergence (primarily because of repeated use of information in data fusion). In practice, we have to contend with a compromise between fast convergence, accurate information sharing and low storage and communication cost. Inspired by these findings, a problem-oriented hybrid protocol that takes the advantages of different approaches while minimizing the side-effects will be valuable.

VI. CONCLUSION

Flooding is an efficient albeit simple solution for information sharing over networks and is the basis of many other networking protocols. In this paper, we formulated it from a set-theoretic perspective, named distributed set-theoretic information flooding (DSIF). This led to a novel consensus protocol for networking referred to as collecting consensus, which has significant both advantages and disadvantages over averaging consensus and diffusion. We have analyzed the explicit convergence and optimality of DSIF based on a novel metric of DoC (degree of consensus). Practical solutions have been proposed either to determine the minimum number of iterations required for any desired DoC or to calculate the DoC that can be achieved by an actual number of iterations. It has also been noted that to save communication, data fusion (such as averaging) can be employed during flooding, which however may cause repeated information use and slower convergence. This trade-off has been analyzed.

Based on the theoretical results, a distributed particle filter framework is proposed and implemented for nonlinear target tracking which applies DSIF on sensor data alone or jointly with intermediate estimates. Simulations have demonstrated the convergence of the DSIF (faster than averaging consensus), the relationship between the filter performance and the DoC, and the advantage and disadvantage of applying parameterized approximation and data fusion for networking.

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as the likelihood consensus code from Dr. O. Hlinka. The authors would also like to acknowledge one anonymous reviewer's idea for a protocol which 1) applies DSIF for a certain number of time intervals, then 2) reduces the amount of stored information through a weighted averaging — like step, and 3) restarts with DSIF approach. We believe that this is a valuable research topic.

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